

# Thinking Abstractly About Constraint Modelling II

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# Previously on the X-files...

- When viewed **abstractly**, many combinatorial problems that we wish to tackle with constraint solving exhibit **common features**.
- By recognising these commonly-occurring patterns, and
- Developing corresponding **modelling patterns** for representing and constraining these combinatorial objects,
- We can **reduce effort** required when modelling a new problem.

# Previously on the X-files...

- We saw a number of individual patterns:
  - Sequences.
  - (Multi-)Sets.
  - Relations.
  - Functions.

# Previously on the X-files...

- We saw how modelling can introduce **equivalence classes** of assignments

KisSequence	1	2	3	4	5	6
	1	2	0	0	0	0

KisSequence	1	2	3	4	5	6
	1	0	0	2	0	0

- Need to be aware of this happening, know how to counter it.
- Reduces the need for detection of such equivalences.

# In This Episode

- We will see how these individual patterns can be **combined** to model more complex problems.

# Nesting

# Nesting Overview

- We've seen how to model several combinatorial objects.
- Often, problems require us to find one combinatorial object **nested** inside another.
  - A set of sets,
  - A sequence of functions...

# How Common Are Problems Involving Nesting?

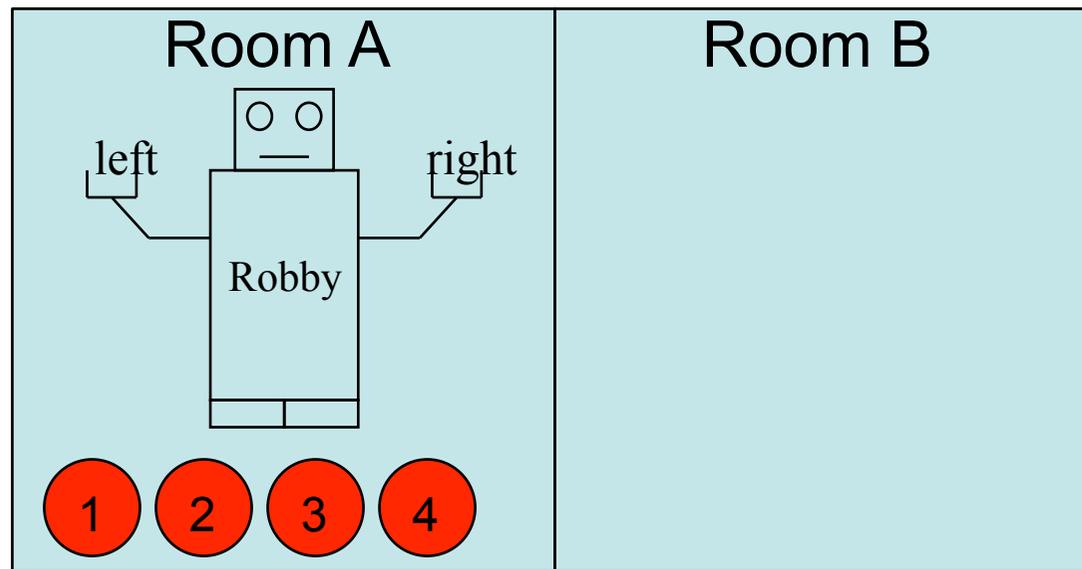
- Very.
- Recall the Steiner Triple (CSPLib 44) problem:
  - Given  $n$ , find a set of  $n(n-1)/6$  triples of elements from  $1, \dots, n$  such that any pair of triples have at most one common element.
  - If  $n = 7$ :
    - $\{\{1, 2, 3\}, \{1, 4, 5\}, \{1, 6, 7\}, \{2, 4, 6\}, \{2, 5, 7\}, \{3, 4, 7\}, \{3, 5, 6\}\}$
  - This is a **set of sets** (the triples).

# How Common Are Problems Involving Nesting?

- Planning Problems:
  - Find a **sequence of actions** to transform an initial state into a goal state.
- When a planning problem allows us actions to be performed in parallel in a single step, it is natural to characterise it as a **sequence of sets** of actions.

# How Common Are Problems Involving Nesting?

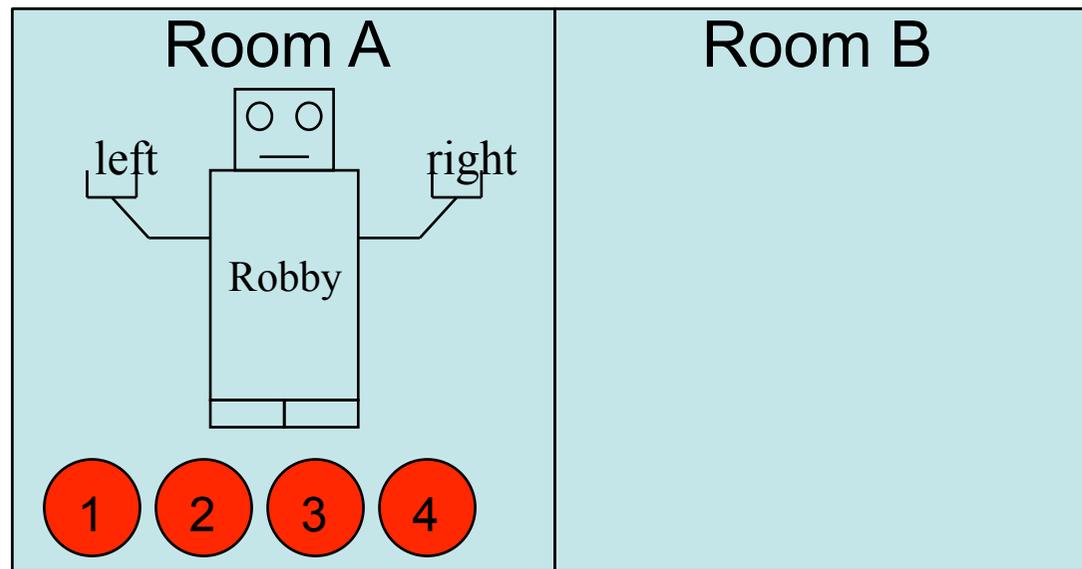
- Example: The Gripper Problem



- **Goal:** All balls in Room B.
- **Operators:** pick up, put down, move.

# How Common Are Problems Involving Nesting?

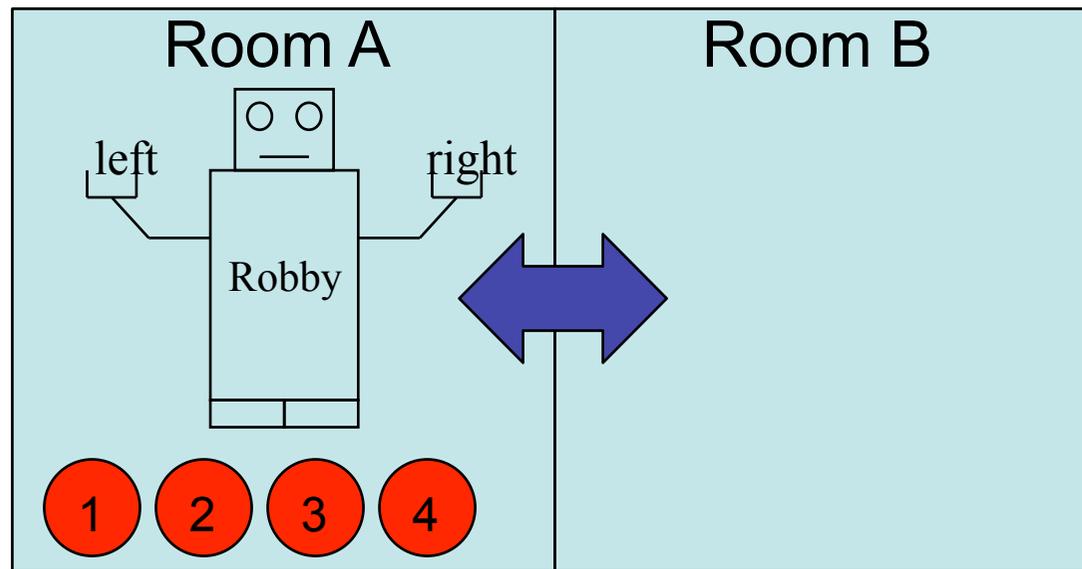
- Example: The Gripper Problem



- Since Robby has two grippers, in a single step of the plan he can pick up/put down two balls.

# How Common Are Problems Involving Nesting?

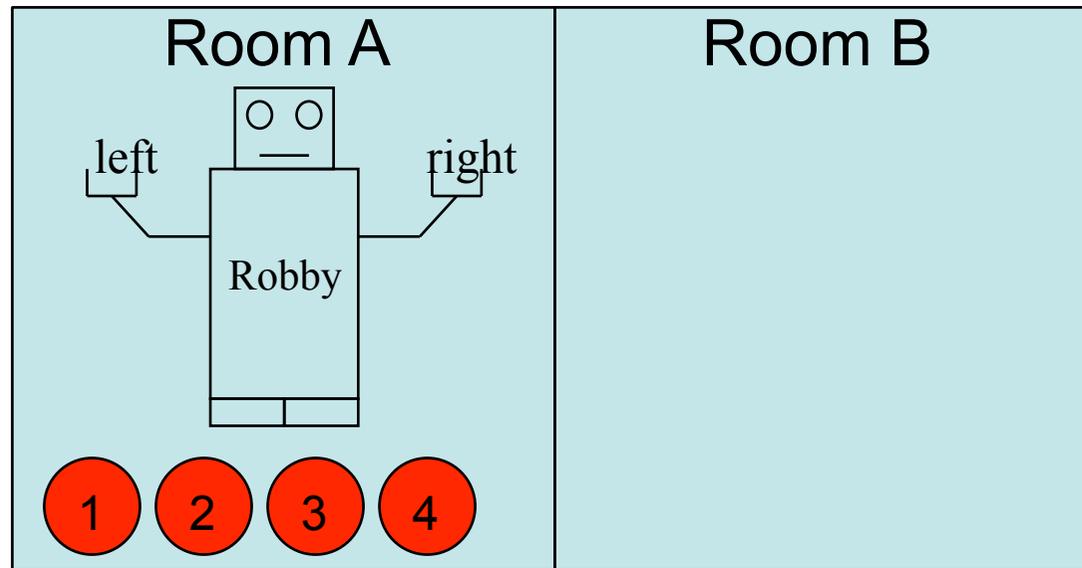
- Example: The Gripper Problem



- When Robby moves, he can't pick up/put down.
- So at most **2** actions per step.

# How Common Are Problems Involving Nesting?

- Example: The Gripper Problem



- Natural to characterise this problem as finding a **sequence of sets** of maximum cardinality 2.

# How Common Are Problems Involving Nesting?

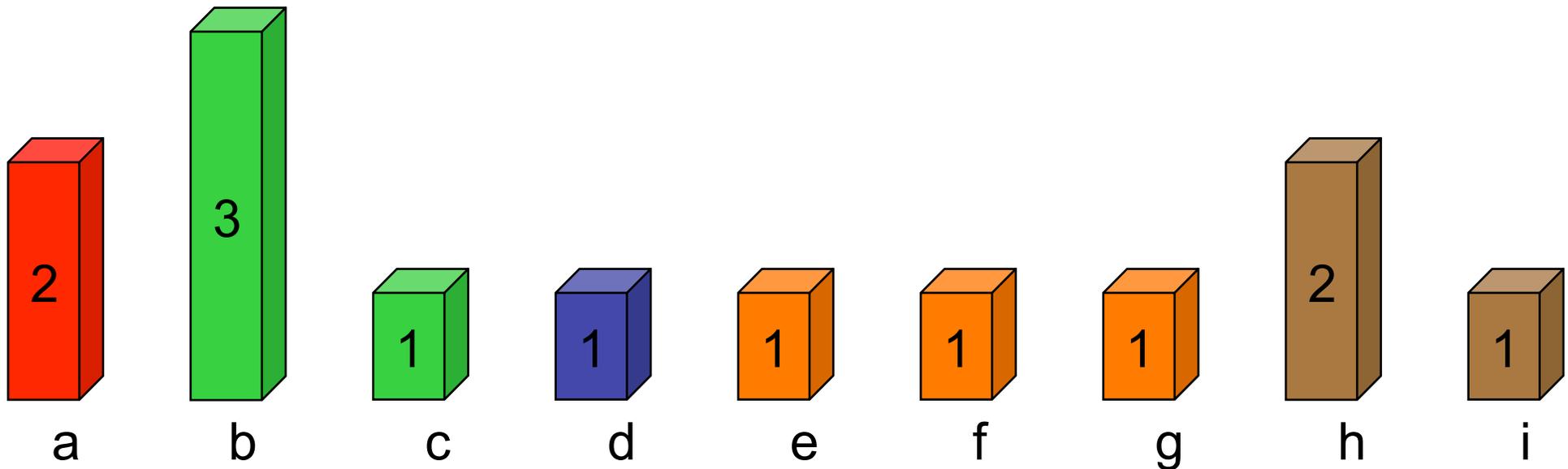
- Example: Steel Mill Slab Design (CSPLib 38).
- The mill can make  $\sigma$  different slab sizes.
- Given  $d$  input orders with:
  - A *colour* (route through the mill).
  - A *weight*.
- Pack orders onto slabs such that the total slab capacity is minimised, subject to:
  - Capacity constraints.
  - Colour constraints.

# How Common Are Problems Involving Nesting?

- Example: Steel Mill Slab Design (CSPLib 38).
- **Capacity:**
  - Total weight of orders assigned to a slab cannot exceed slab capacity.
- **Colour:**
  - Each slab can contain at most  $p$  of  $k$  total colours.
  - Reason: expensive to cut slabs up to send them to different parts of the mill.

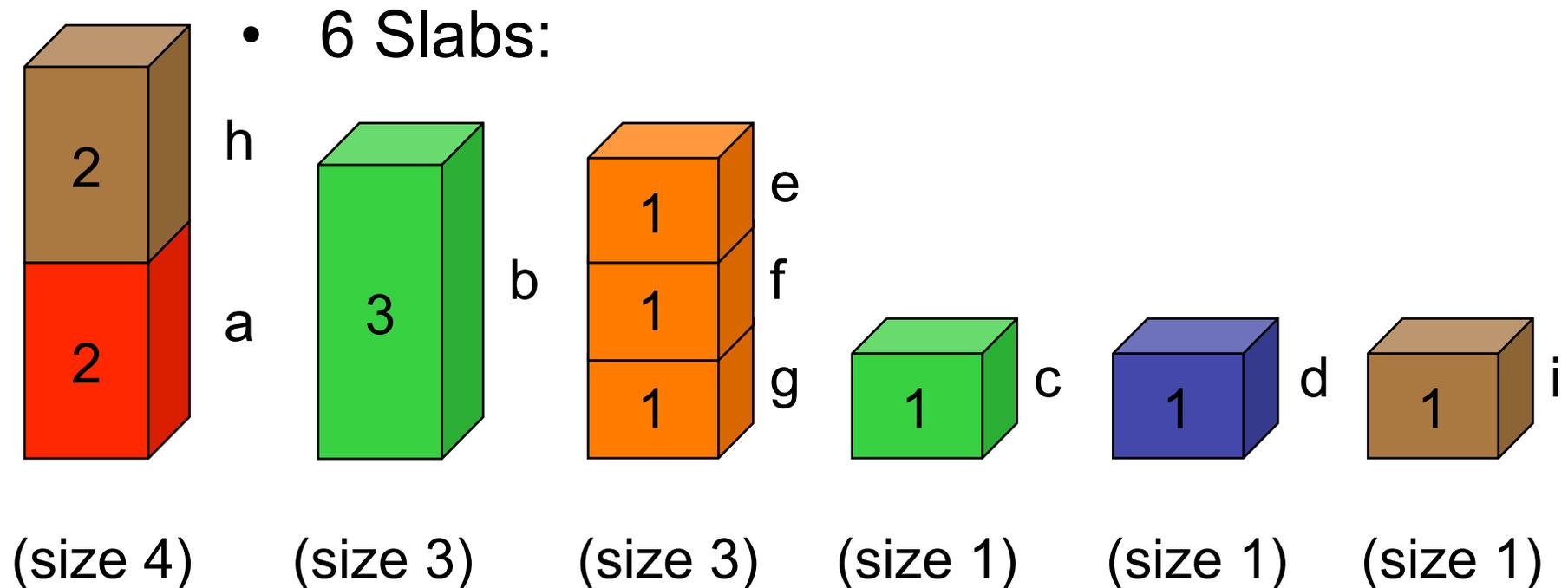
# How Common Are Problems Involving Nesting?

- Example: Steel Mill Slab Design (CSPLib 38).
- Slab Sizes:  $\{1, 3, 4\}$  ( $\sigma = 3$ )
- Orders:  $\{o_a, \dots, o_i\}$  ( $d = 9$ )
- Colours:  $\{\text{red, green, blue, orange, brown}\}$  ( $k = 5$ )
- $p = 2$



# How Common Are Problems Involving Nesting?

- Example: Steel Mill Slab Design (CSPLib 38).
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# How Common Are Problems Involving Nesting?

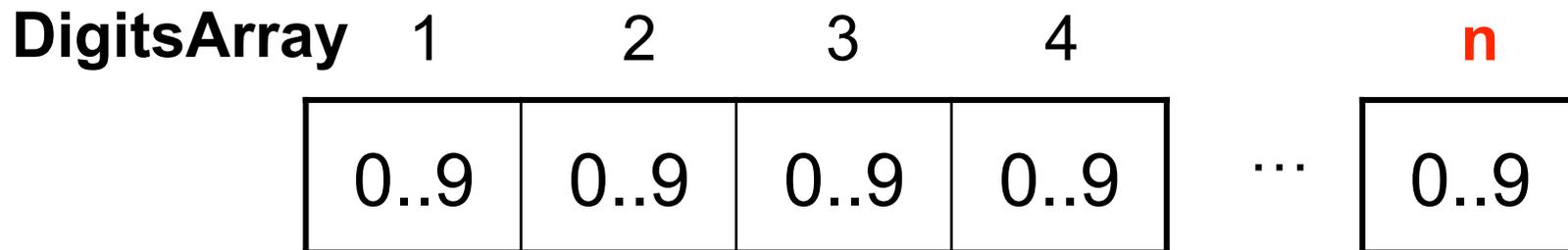
- Example: Steel Mill Slab Design (CSPLib 38).
- A slab can be represented as a **set** of orders.
- We must also determine the size of each slab.
- So this problem can be characterised as a **function from sets of orders to the set of sizes**.
- The function is **partial**, since not all possible sets of orders will be mapped to a slab size.

The Template design problem (CSPLib 2) can be characterised similarly.

# **Nesting Inside Sequences**

# Nesting Inside Sequences

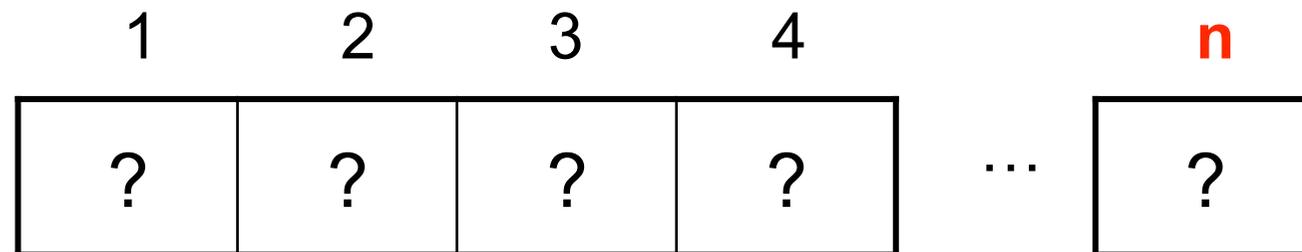
- Recall how we modelled fixed-length sequences.
- An array of decision variables indexed  $1..n$ . Domains are the objects to be found.
- Example, find a sequence of  $n$  digits:



# Nesting Inside Sequences

- Assume now that we must find a sequence of sets, functions, relations, ...
- We can no longer use a single variable at each index to represent the object at that position.
  - Because 1 variable is not enough to represent our set, function or relation.

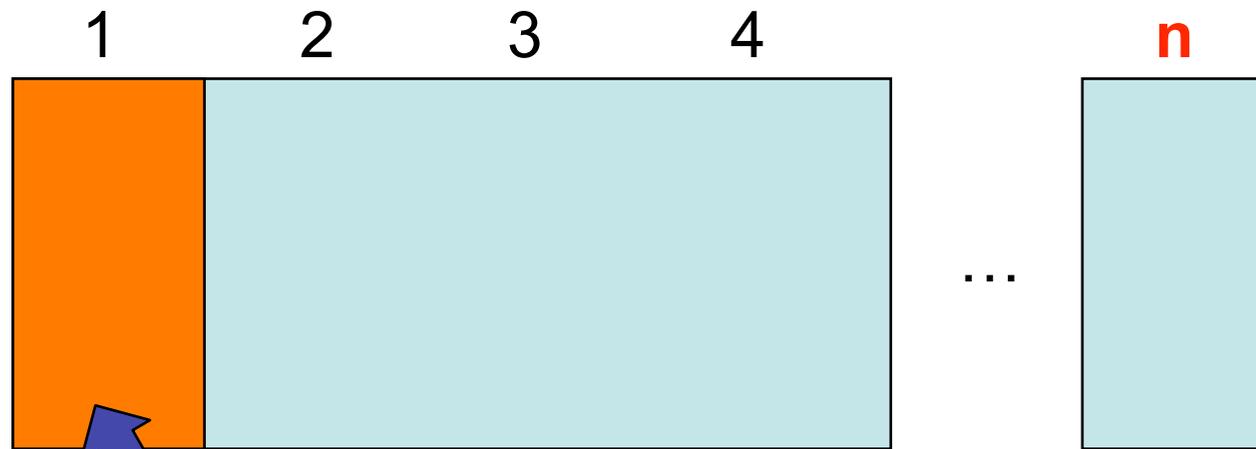
## NestedSeqArray



# Nesting Inside Sequences

- Simple solution:
  - Extend the dimension of the array.

## NestedSeqArray

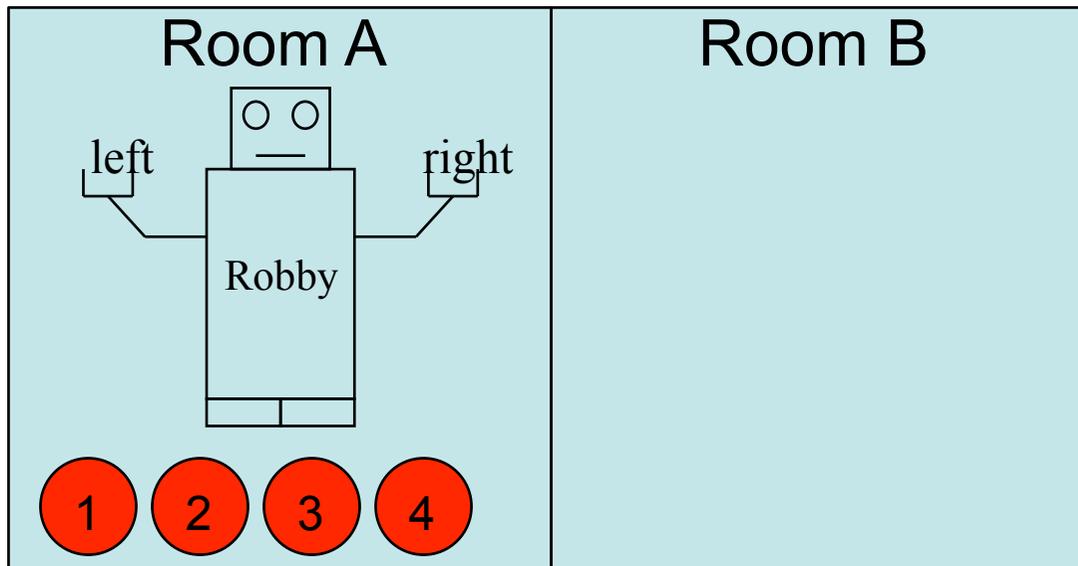


E.g. 2 dimensions.

We now have a **column** of variables to represent our set, relation, function ...

# Nesting Inside Sequences

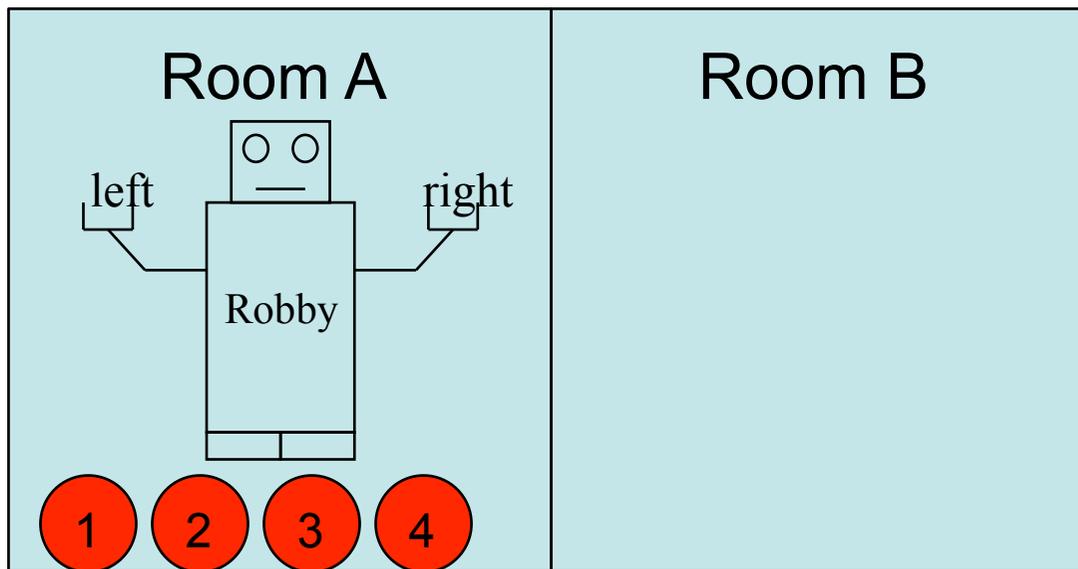
- To illustrate, consider modelling a sequence of sets.



- Returning to the Gripper problem, assume that we are looking for a plan of length **n**.

# Nesting Inside Sequences

- We saw before that a set of cardinality at most two can be used to model the actions performed at each step.



- Need elements for moving, pick up/drop balls with each of the two grippers.
- Assume we use integers  $1..k$  to represent these actions.
- (I'm glossing over details here).

# Nesting Inside Sequences

- So, we have a sequence of length **n** of sets of cardinality at most 2 drawn from  $1..k$ .
- Let's start by looking at the **occurrence** representation:

Each column represents the occurrence representation of a set.

Constraints?

	1	2	3	4	...	<b>n</b>
1	0, 1	0, 1	0, 1	0, 1	...	0, 1
2	0, 1	0, 1	0, 1	0, 1	...	0, 1
...						...
<b>k</b>	0, 1	0, 1	0, 1	0, 1	...	0, 1

# Nesting Inside Sequences

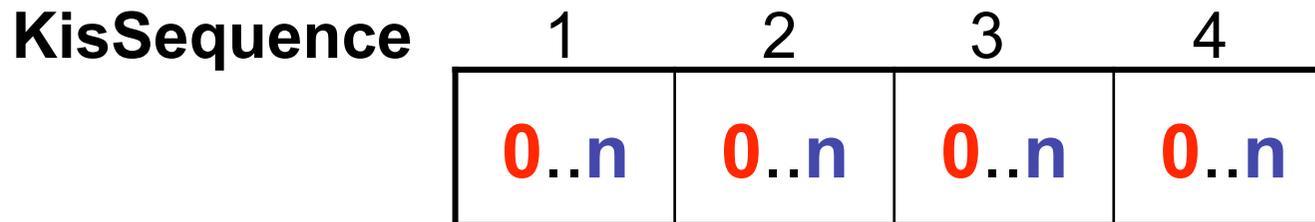
- So, we have a sequence of length **n** of sets of cardinality at most 2 drawn from  $1..k$ .
- Now let's look at the **explicit** representation.

Each column represents the explicit representation of a set.  
Constraints?

	1	2	3	4	...	<b>n</b>
1	0..k	0..k	0..k	0..k	...	0..k
2	0..k	0..k	0..k	0..k		0..k

# Nesting Inside Sequences

- What if the sequence has bounded length?
- Recall that in the non-nested case we used a dummy value:



# Nesting Inside Sequences

- We can use the same approach here (careful not to use the same dummy value as the explicit model of the inner sets).
- Could also use auxiliary **switch** variables to indicate whether the corresponding column is part of the sequence.
- Again, careful of introducing equivalence classes of assignments.

	1	2	3	4	...	<b>n</b>
1	0..k	0..k	0..k	0..k	...	0..k
2	0..k	0..k	0..k	0..k	...	0..k
Switches	0, 1	0, 1	0, 1	0, 1	...	0, 1

# Nesting Inside Sets

# Nesting Inside Sets

- Being asked to find a set of some other object is common, so it is worth considering how to model this type of problem.
- Now we must choose how to model the outer type (e.g. explicit vs occurrence model of sets) **as well** as the inner.

# Nested Sets

Consider the following simple problem class:

- Given **m**, **n**.
- Find a cardinality-**m** **set of sets** of **n** digits such that ...
- From what we have seen so far, we have three possibilities:
  1. An occurrence representation.
  2. Outer: Explicit. Inner: Occurrence.
  3. Outer: Explicit. Inner: Explicit.

# Nesting Inside Sets: Occurrence

- Recall the occurrence representation of a fixed-cardinality set of digits:

○

0	1	2	3	4	5	6	7	8	9
0, 1	0, 1	0, 1	0, 1	0, 1	0, 1	0, 1	0, 1	0, 1	0, 1

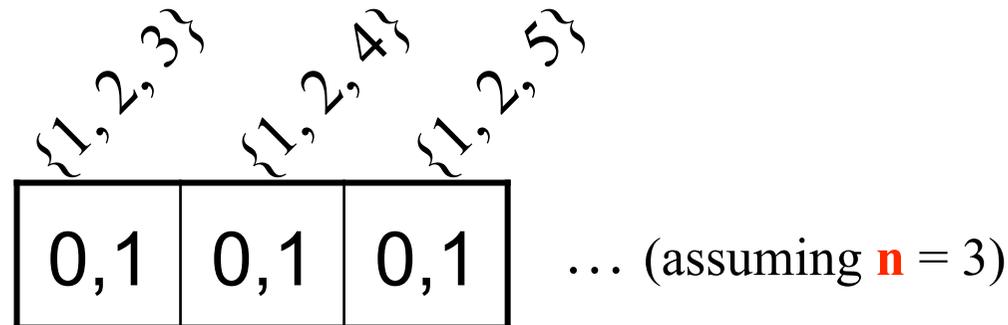
- We have an index per possible element of the set.

# Nesting Inside Sets: Occurrence

Can we take the same approach here?

- Given **m**, **n**.
- Find a cardinality-**m** set of sets of **n** digits such that...

Introduce an array indexed by the possible sets of **n** digits!

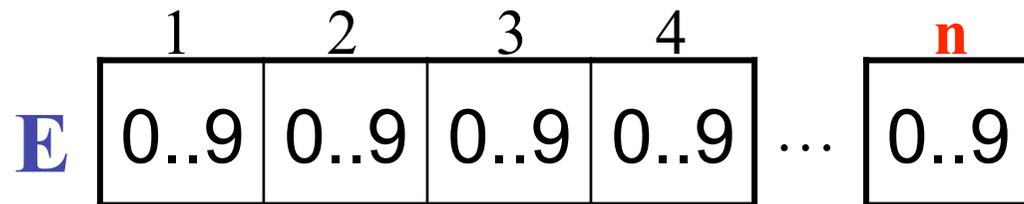


This is often not feasible.

Typically, when dealing with nesting the outer layers are represented **explicitly**.

# Nesting Inside Sets: Outer Explicit

- Recall the explicit representation of a fixed-cardinality set of digits:



- Similarly to the sequence example, we extend the dimension of **E** according to the representation we choose for the inner set.
- We're also going to have to be careful to make sure the elements of the outer set are **distinct**.

# Nesting Inside Sets: Explicit/Occurrence

- Given **m**, **n**.
- Find a cardinality-**m** set of sets of **n** digits such that...
- Let's consider an occurrence representation for the inner sets.

But what about **equivalence classes?**

<b>EO</b>	1	2	...	<b>m</b>
0	0,1	0,1		0,1
1	0,1	0,1	...	0,1
2	0,1	0,1		0,1
		...		...
9	0,1	0,1	...	0,1

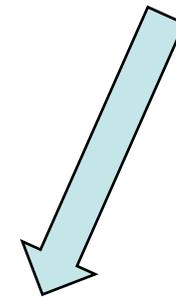
Constraints:

$$\text{Sum}(\text{col } i \text{ of } \mathbf{EO}) = n$$

(foreach  $i$  in  $1..m$ )

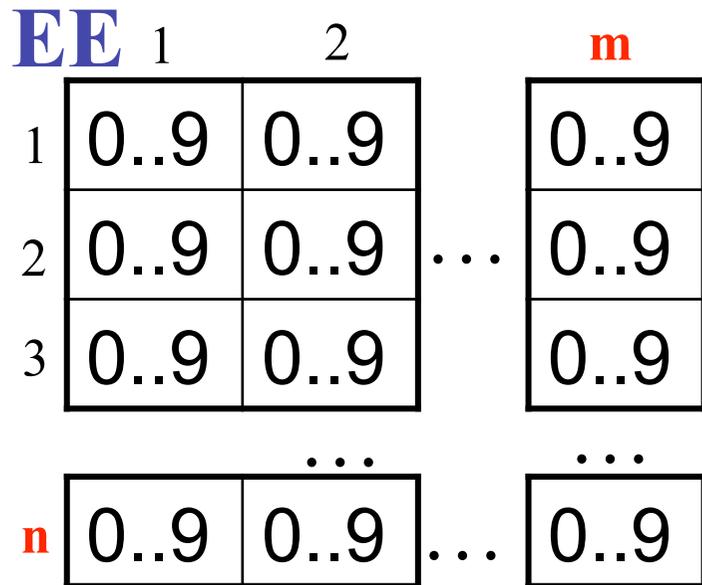
$$\text{Scalar-prod}(\text{col } i \text{ of } \mathbf{EO}, \text{col } j \text{ of } \mathbf{EO}) \neq n$$

(foreach  $\{i, j\}$  in  $1..m$ )



# Nesting Inside Sets: Explicit/Explicit

- Given **m**, **n**.
- Find a cardinality-**m** set of sets of **n** digits such that...
- Let's consider an occurrence representation for the inner sets.



Constraints:

Col  $i$  of **EE**  $<_{\text{lex}}$  Col  $j$  of **EE**  $\vee$

Col  $i$  of **EE**  $>_{\text{lex}}$  Col  $j$  of **EE**

(foreach  $\{i, j\}$  in  $1..m$ )

AllDiff on columns.

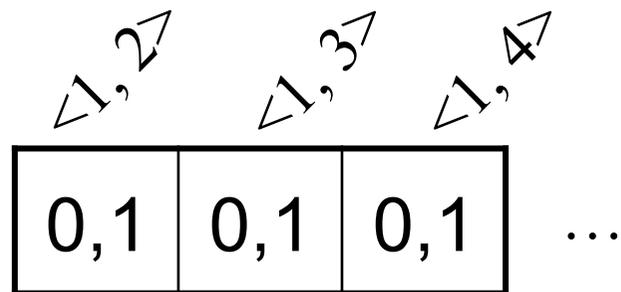
But what about **equivalence classes?**

# Relations as Sets of Tuples

- Last time we looked at a couple of ways of modelling relations.
- We can also view relations as **sets of tuples**.
- Recall our example:
  - Find a relation  $R$  between sets  $A = \{1, 2, 3\}$  and  $B = \{2, 3, 4\}$  such that...
- What happens when we try and model this from the perspective of a set of tuples?

# Relations as Sets of Tuples: Occurrence

- We have an array indexed by the possible tuples:



- Basically same as the occurrence representation we came up with directly:

		A		
		1	2	3
B	2	0, 1	0, 1	0, 1
	3	0, 1	0, 1	0, 1
	4	0, 1	0, 1	0, 1

# Relations as Sets of Tuples: Explicit

- Find a relation R between sets  $A = \{1, 2, 3\}$  and  $B = \{2, 3, 4\}$  such that...
- Maximum number of tuples is 9. Invoke our bounded-cardinality set pattern:

	1	2	3	4	...
1	1..3	1..3	1..3	1..3	
2	2..4	2..4	2..4	2..4	

What about **equivalence classes**?

What if the relation allows fewer than the full 9 tuples?

# **The Social Golfers Problem**

# The Social Golfers Problem

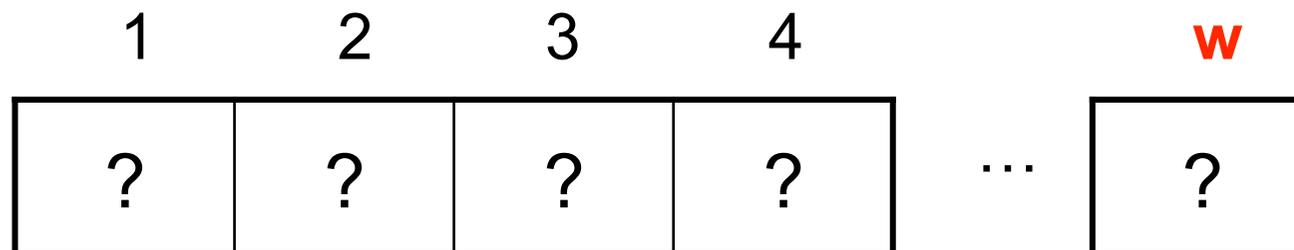
- In a golf club there are a number of golfers who wish to play together in  $g$  groups of size  $s$ .
- Find a schedule of play for  $w$  weeks such that no pair of golfers play together more than once.

# The Social Golfers Problem: Modelling

- In each week, we need to **partition** the golfers into groups.
  - A partition is a set of sets. No pair of inner sets have an element in common.
- What about the weeks?
  - A sequence? But what does the order matter?
  - A multiset.
  - In fact, there's an **implied constraint** here. Can you see it?
- So we can think of the problem as finding a **multiset of partitions**.

# Golfers: Representing the Outer Multiset

- We have seen explicit and occurrence representations of multisets.
- The multiset contains complex objects (partitions).
- Indexing an array by the possible partitions of golfers doesn't seem appealing.
- So let's try an explicit model:

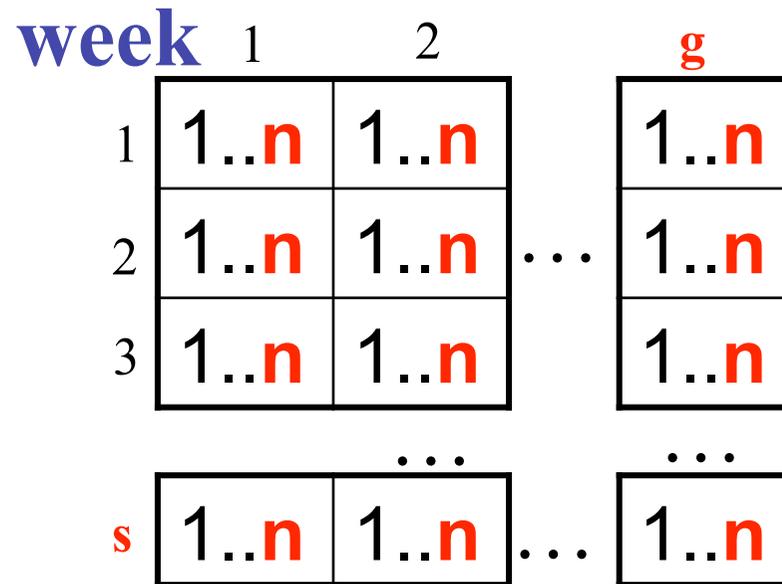


# Golfers: The Partitions

- In each week we want to partition the golfers into **g** groups of size **s**.
- That is, a set of cardinality **g** of sets of cardinality **s**.
- As per the previous discussion, probably sensible to represent the outer set explicitly.
- The inner set could be occurrence or explicit. Here we'll talk about an explicit/explicit representation.

# Golfers: The Partitions

- Let  $n$  = number of golfers =  $g * s$ .

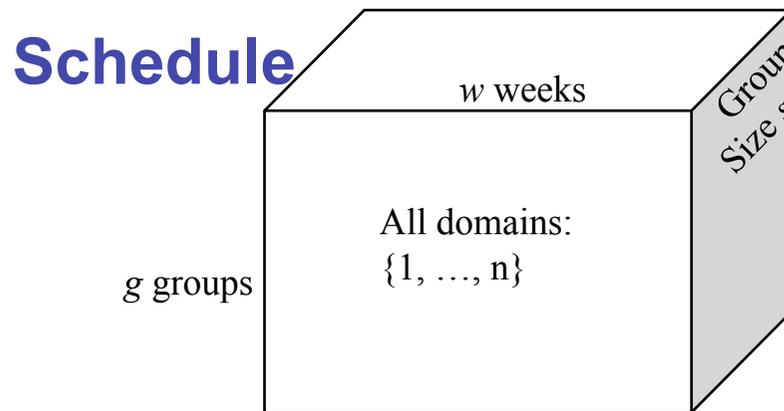


Since a week is a partition, what can we say about the elements of **week**?

What about **equivalence classes**?

# A Multiset of Partitions of Golfers

- If we put **week** into each slot of our multiset representation, we obtain a 3d array:

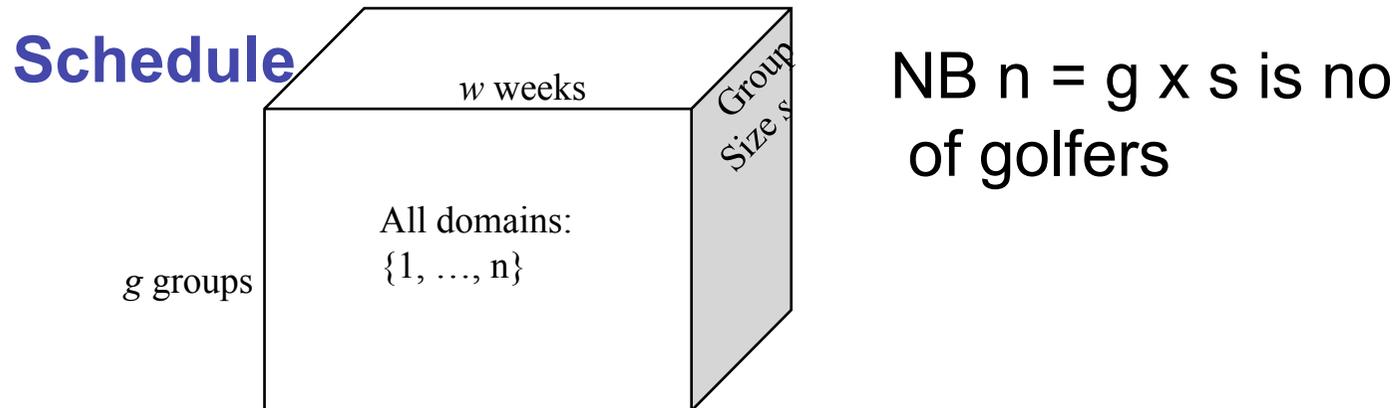


NB  $n = g \times s$  is no of golfers

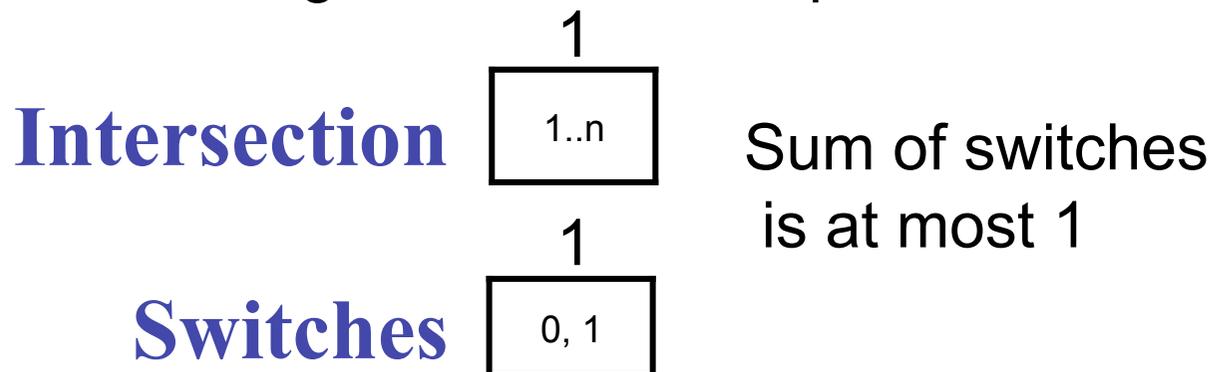
We can **order the weeks** lexicographically to counter the equivalence of assignments obtained by permuting the weeks.

# A Multiset of Partitions of Golfers

- Need to ensure no pair of golfers meet more than once.



Equivalently: size of intersection of each pair of groups is at most 1. Invoking our intersection pattern:



# Social Golfers

- Solution to the instance with 3 groups (size 3) over 3 weeks:

3 groups, size 3

	[1, 2, 3]	[4, 5, 6]	[7, 8, 9]
3 weeks	[1,4,7]	[2,5,8]	[3,6,9]
	[1,5,9]	[2,6,7]	[3,4,8]

We've missed an equivalence class! Can you spot it?  
Hint: we saw something similar in the BIBD.

# Nesting Summary

- Modelling problems involving nested combinatorial objects can be quite tricky.
- Using the patterns we've been looking at can help you to do it **systematically**.
- It can also help in spotting **equivalence classes** of assignments as you introduce them.
  - Which can be substantially cheaper than trying to detect them after the fact.

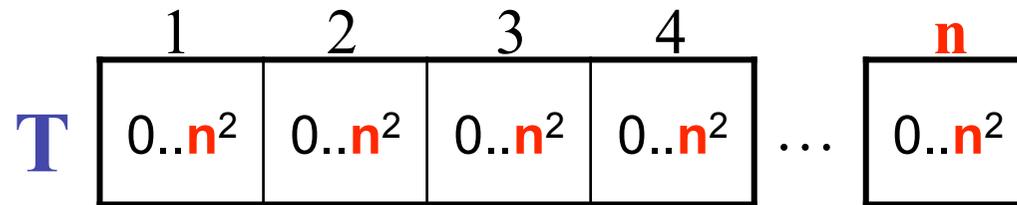
And Finally:

# **The Golomb Ruler Challenge**

# The Golomb Ruler Problem

- NB This is a type of **Graceful Graph**.
- Given:
  - A positive integer  $n$ .
- Find:
  - A set of  $n$  integer ticks on a ruler of length  $m$ .
- Such that:
  - All inter-tick distances are distinct.
- Minimising:
  - $m$ .

# Modelling the Golomb Ruler



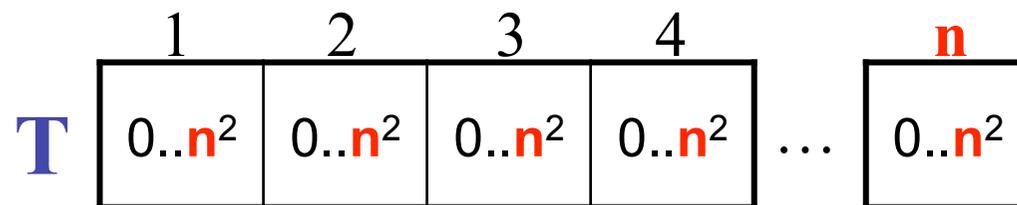
- All inter-tick distances are distinct:
  - $\mathbf{T}[j] - \mathbf{T}[i] \neq \mathbf{T}[k] - \mathbf{T}[l]$   
for each  $\{i, j\}, \{k, l\}$  drawn from  $1..n$ ,  
such that  $\{i, j\} \neq \{k, l\}, i < j, k < l$   
again, exploiting ascending order.
- Objective:
  - Minimise( $\mathbf{T}[n]$ )  
Again, exploiting ascending order.

# Modelling the Golomb Ruler

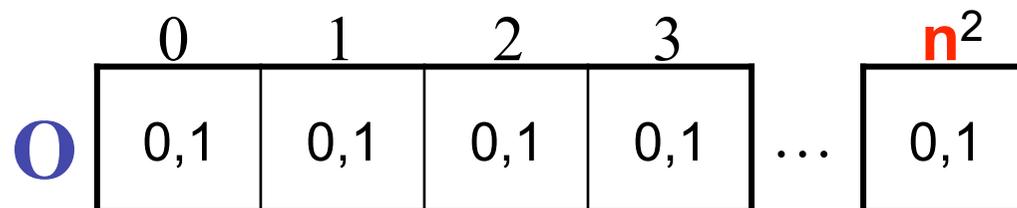
- **A Challenge:**
- Can you see how to model this problem using the occurrence representation?
- This does require a little sleight of hand...

# Golomb Ruler: Occurrence Model

- Recall that in our explicit model, the elements of the set are  $0..n^2$ .

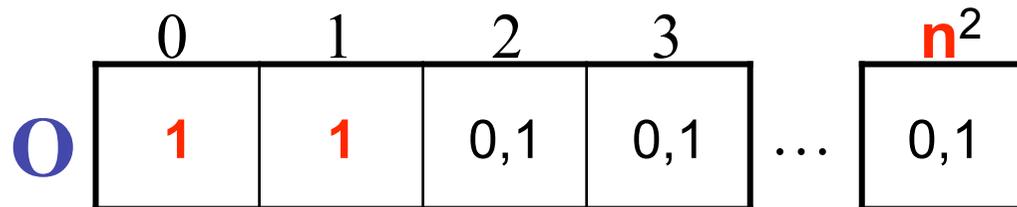


- Invoking our occurrence representation pattern, we begin with an array **O** indexed  $0..n^2$ :



# Golomb Ruler: Occurrence Model

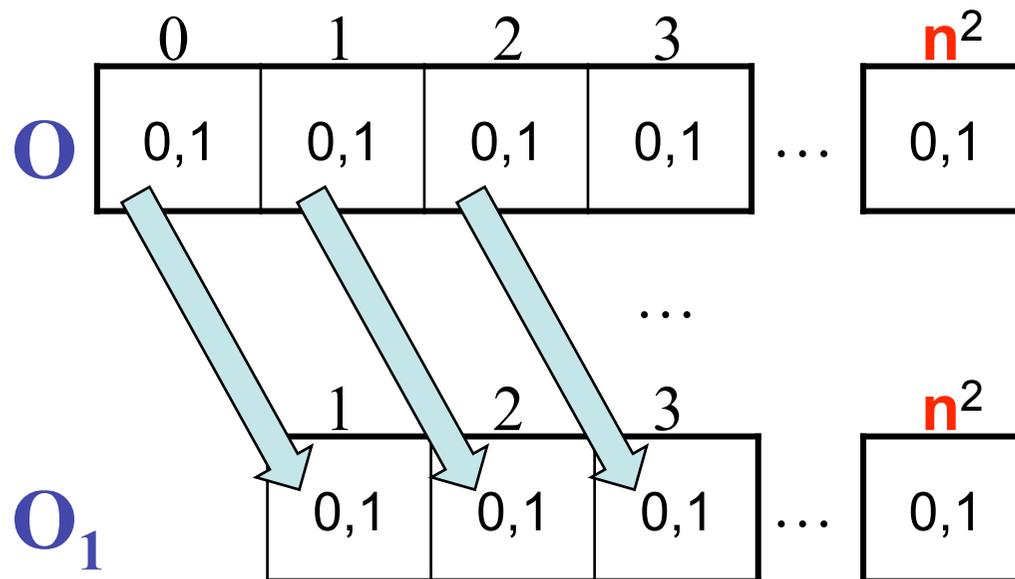
- How can we express the distinct distances constraint?
- Consider a partial assignment:



- We now know that no other pair of adjacent variables can be assigned 1.
- How can we express these constraints?

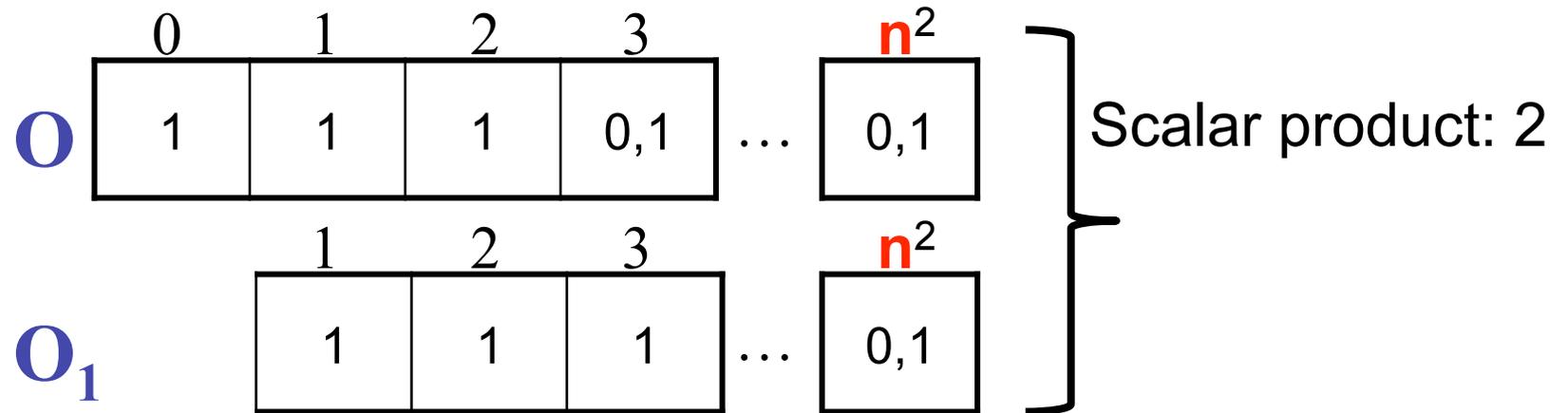
# Golomb Ruler: Occurrence Model

- Consider an array  $O_1$ , which contains the **same** variables as  $O$ , shifted one position right.

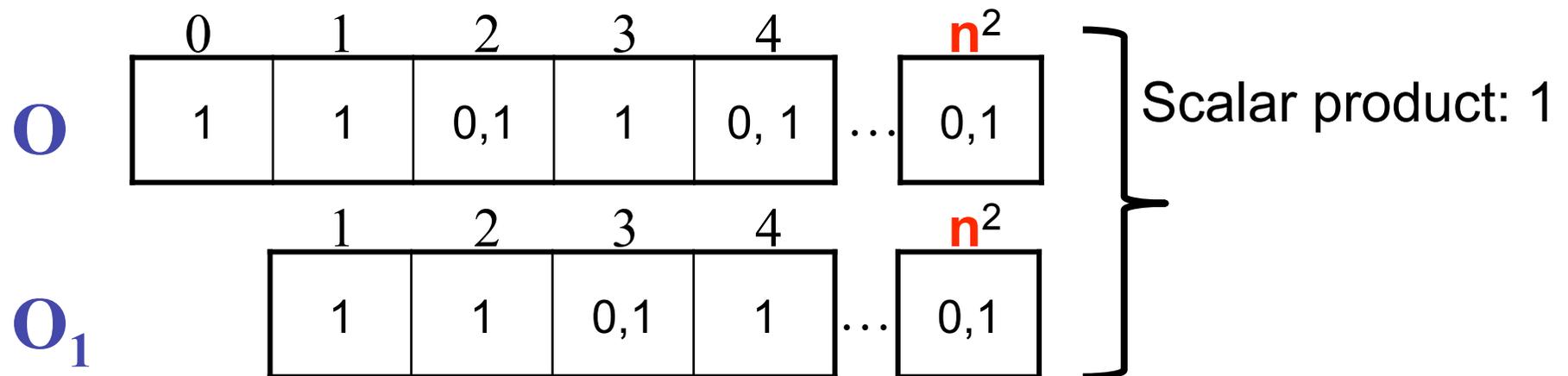


# Golomb Ruler: Occurrence Model

- Now let's assign some variables:

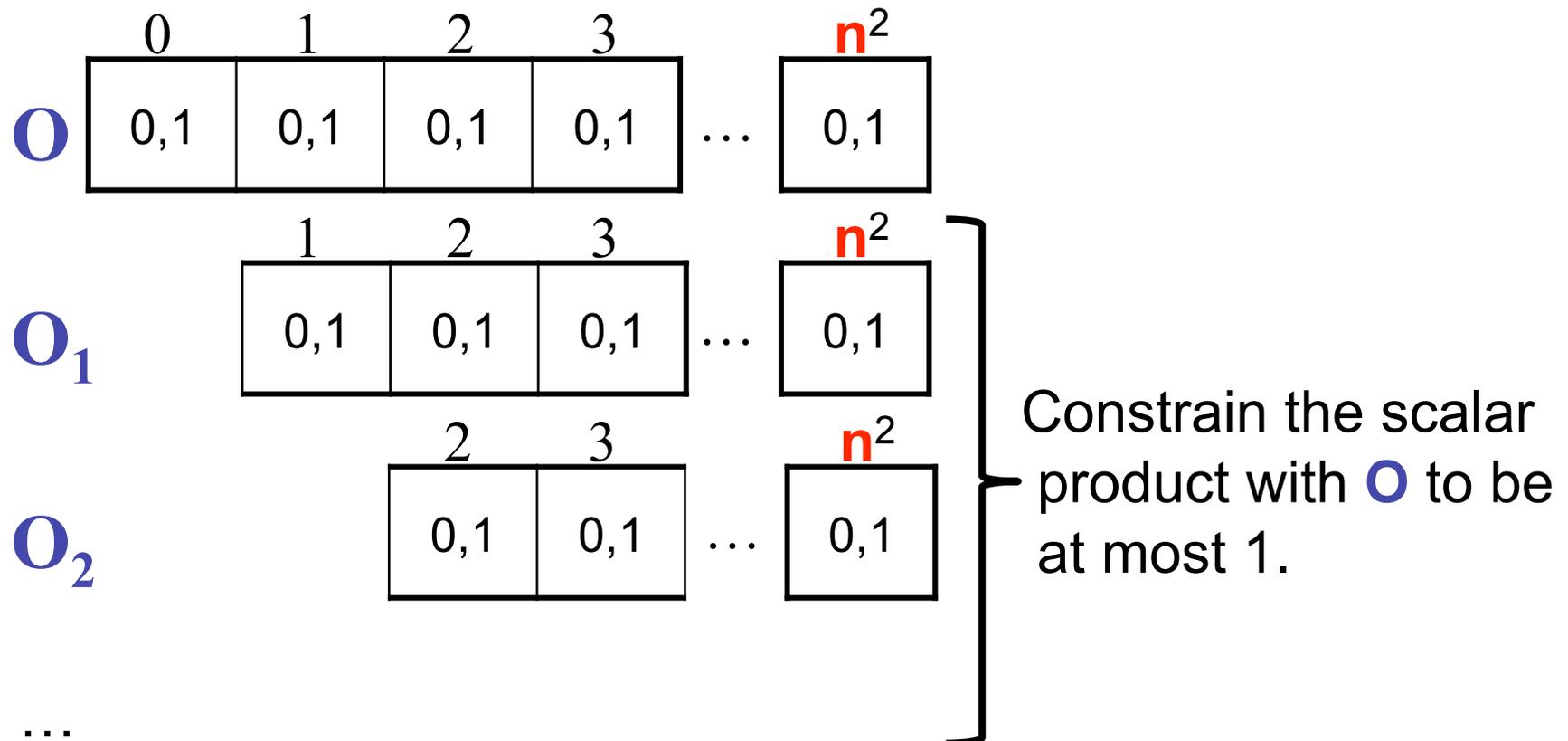


- Whereas:



# Golomb Ruler: Occurrence Model

- Now consider adding one such array per difference:

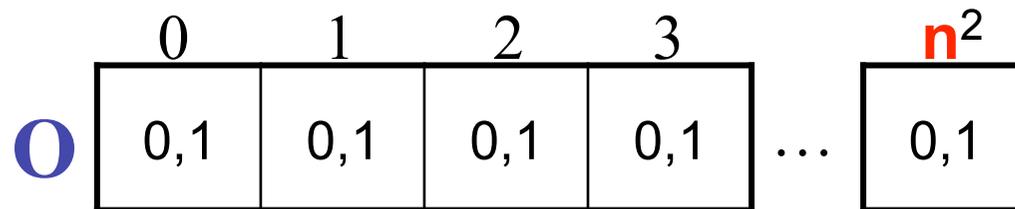


# Golomb Ruler: Occurrence Model

- (Perhaps) you're thinking:
  - That's a lot of extra variables!
- In fact, we've introduced **no** extra variables.
- Just re-used **existing** variables in new arrays.

# Golomb Ruler: Occurrence Model

- But what about the objective?



- Tricky because we are trying to minimise the index of the last “1” assignment.
- One way is to solve a series of problems, increasing the size of  $O$ .
  - As soon as we have a solution, it is optimal.

# Golomb Ruler: Discussion

- Most constraint models of this problem for the literature focus on the explicit representation of the set.
- Build on this model by adding **auxiliary variables** and **implied constraints**.
  - Barbara M. Smith, Kostas Stergiou, Toby Walsh: Using Auxiliary Variables and Implied Constraints to Model Non-Binary Problems. AAI/IAAI 2000: 182-187
- Distributed effort to find large GRs looks more like this occurrence model.
  - The power of bit-shifting.