

# **Constraints and Automata (time-series constraints)**

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- **Background**
- Synthesizing automata with accumulators from transducers
- Parametric glue matrices
- Simplifying automata with accumulators
- Reformulating in LP
- Deriving necessary conditions (as linear constraints)
- Bounds

# The story to remember

Generate a few hundred time-series constraints  
in a **compositional way**,  
but **need**  
to **simplify** them,  
to **generate necessary conditions** and **bounds**  
to get some reasonable behaviour both in CP and LP

Since **too many constraints** this cannot be done  
for each constraint individually in a reasonable time-frame

**Has also to handle the combinatorial aspect  
of these constraints in a compositional way**

# Finite state transducer

(*introduced by* M.P. Schützenberger)

- A FST is defined by:
  - A finite set of **states**  $Q$
  - An **input alphabet**  $\Sigma$  (*finite set of input symbols*)
  - An **output alphabet**  $\Gamma$  (*finite set of output symbols*)
  - A **transition relation**  $\delta: Q \times (\Sigma \cup \{\varepsilon\}) \times (\Sigma \cup \{\varepsilon\}) \times Q$
  - An **initial state**  $q_{\text{init}} \in Q$
  - A set of **accepting states**  $F \subseteq Q$

Some time can have more than one output symbol  
in the transition relation

# Finite state transducer

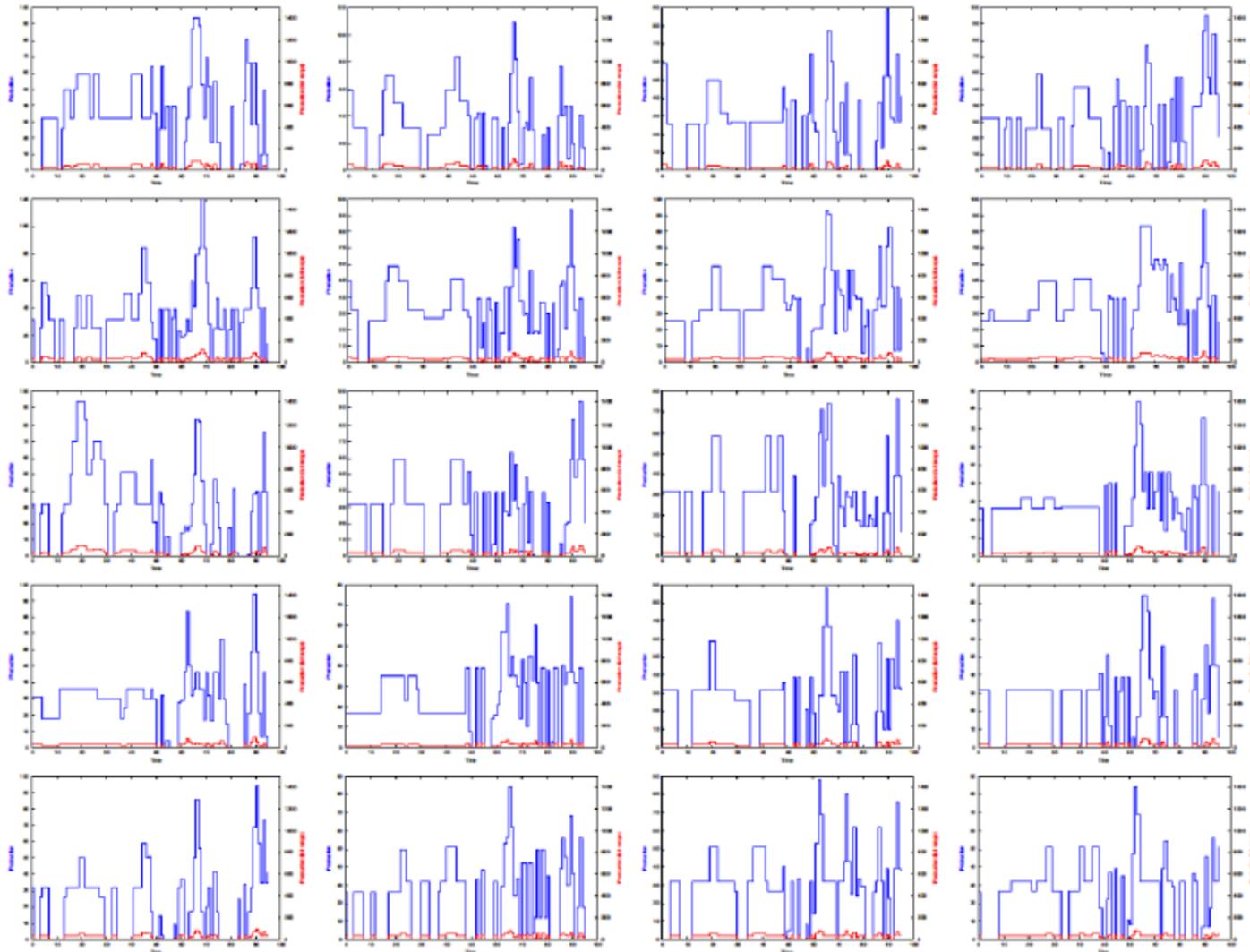
- Used mostly
  - Natural language (*text of speech*)
- But also
  - Computational biology
  - Fraud detection
- Popular (*like automata*) in industry
  - microsoft research
  - google (M. Mohri)
- Like automata you can learn them

examples of transducers later on

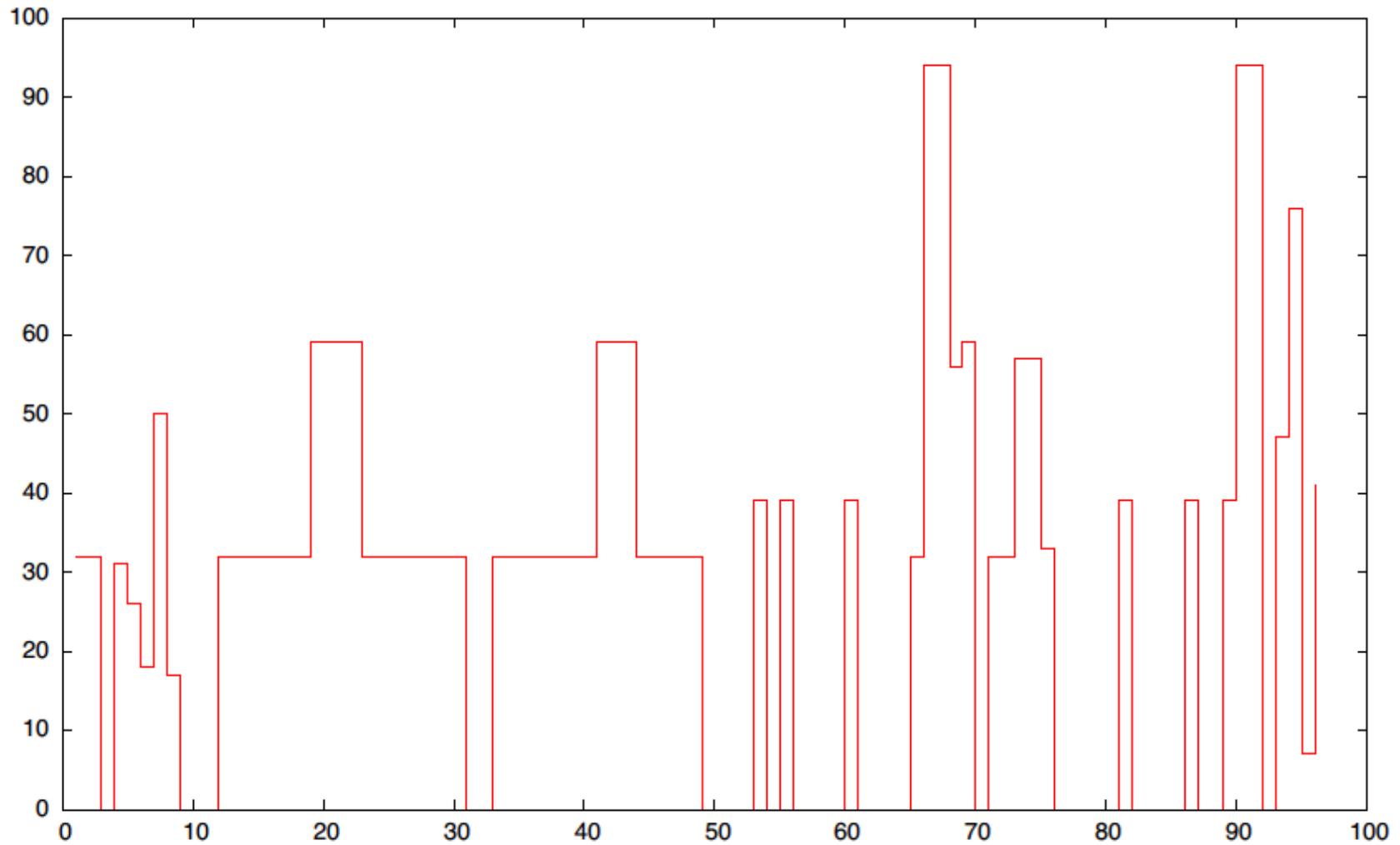
# Background

- PGMO Project with EDF
- Analysis of power output curves for electricity generators
- Use ModelSeeker to describe/categorize/synthesize output from UCP model
- Published in CP 2013

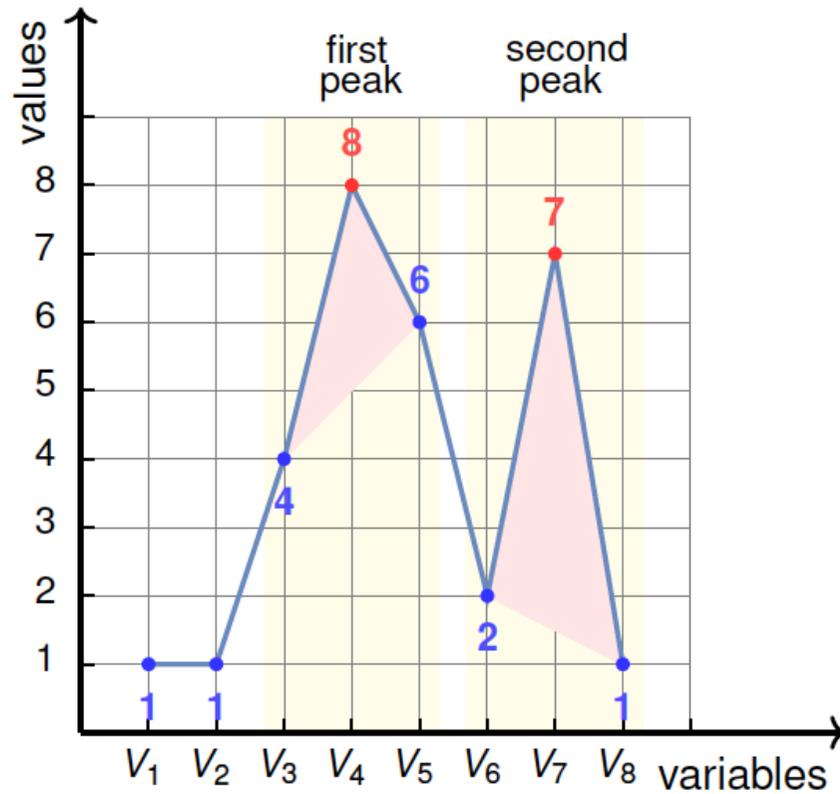
# Example: From this ...



... to this (generated profile)



# Example: the peak constraint



peak  $(2, \langle 1, 1, 4, 8, 6, 2, 7, 1 \rangle)$

# Automaton with counters: peak constraint in Global Constraint Catalog

## STATE SEMANTICS

$s$  : stationary/decreasing mode ( $\{> | =\}^*$ )  
 $u$  : increasing mode ( $\{< | =\}^*$ )

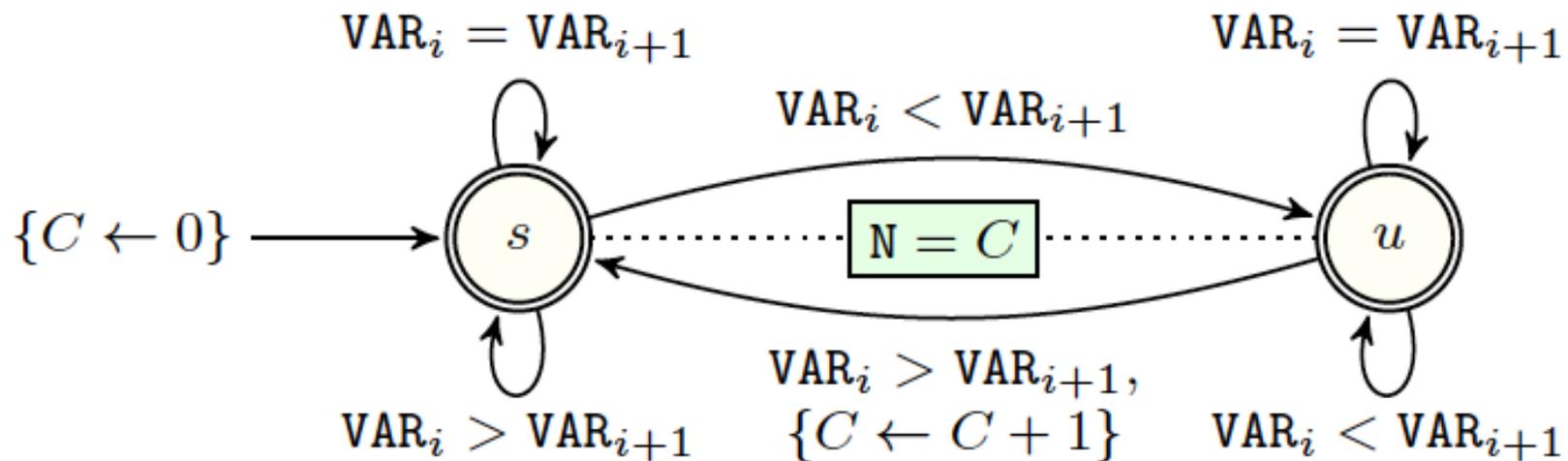


Figure 5.689: Automaton of the PEAK constraint

But always some missing constraints  
*(when meeting people from industry)*

and dont want to introduce the missing  
constraints one by one in the  
global constraint catalog

which leads to a synthesized  
time-series catalog

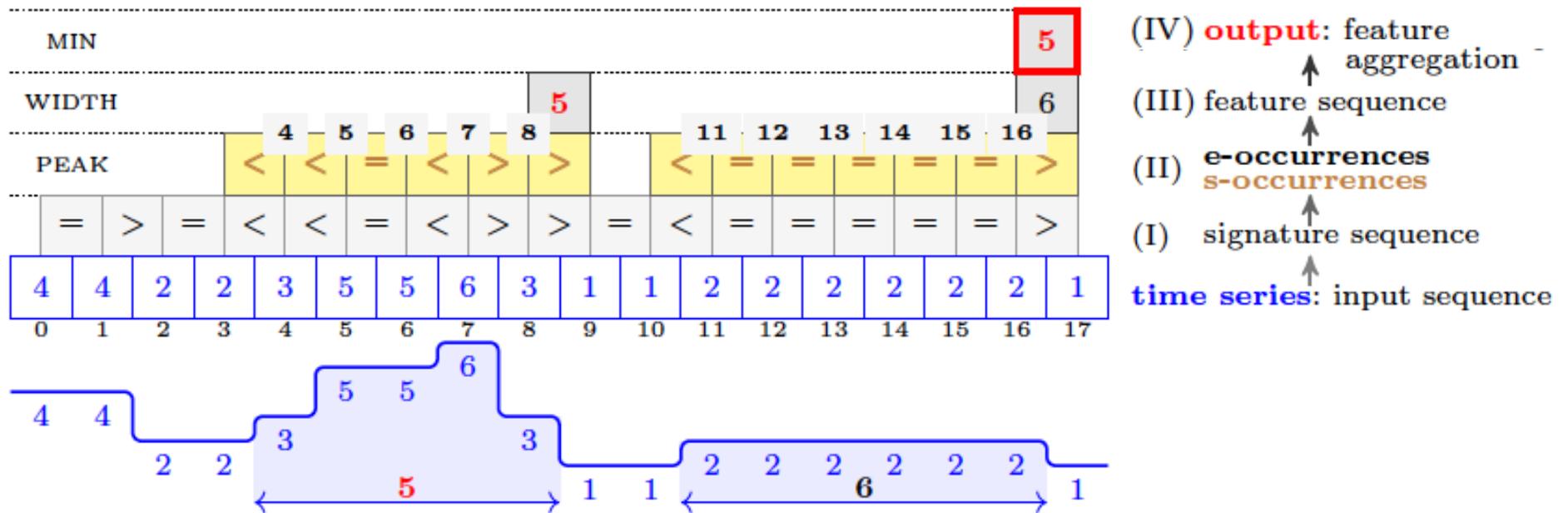
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# Decomposing the definition of a constraint

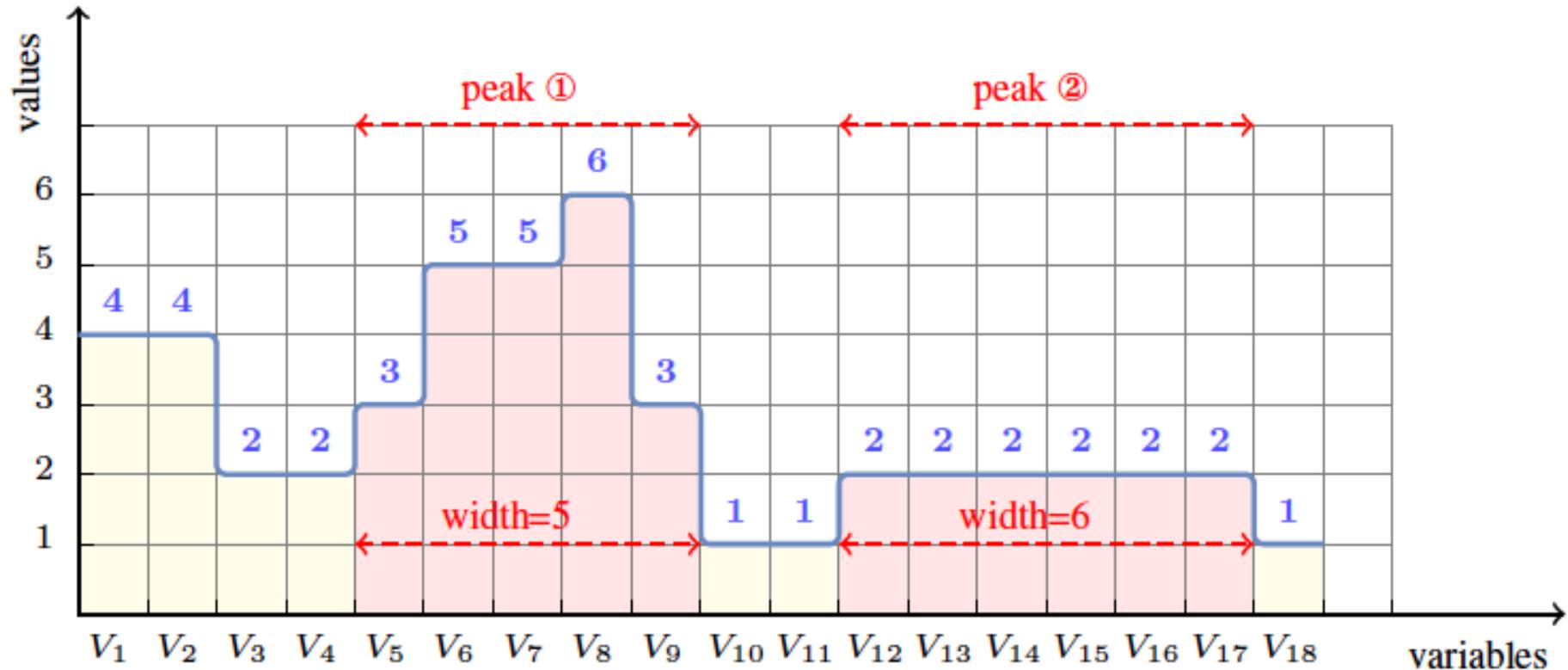
- Constraints are **pure functional dependencies** or **predicates** (*see predicates later on*)
- Implemented as **automata with counters**
- **Four steps** (layers) in definition
  - Building **signature**
  - Recognize **pattern occurrences** in sequence
  - Extract **feature** per pattern
  - **Aggregate** features

# Example: min\_width\_peak



`MIN_WIDTH_PEAK(5, <4, 4, 2, 2, 3, 5, 5, 6, 3, 1, 1, 2, 2, 2, 2, 2, 2, 1>)`

# Example: min\_width\_peak (continued)



`MIN_WIDTH_PEAK(5, <4, 4, 2, 2, 3, 5, 5, 6, 3, 1, 1, 2, 2, 2, 2, 2, 1>)`

# Signature

- Convert (integer) value to finite alphabet
- Signature links two consecutive entries in time-series
- We use  $<, =, >$  with their natural semantics
- Other signatures possible

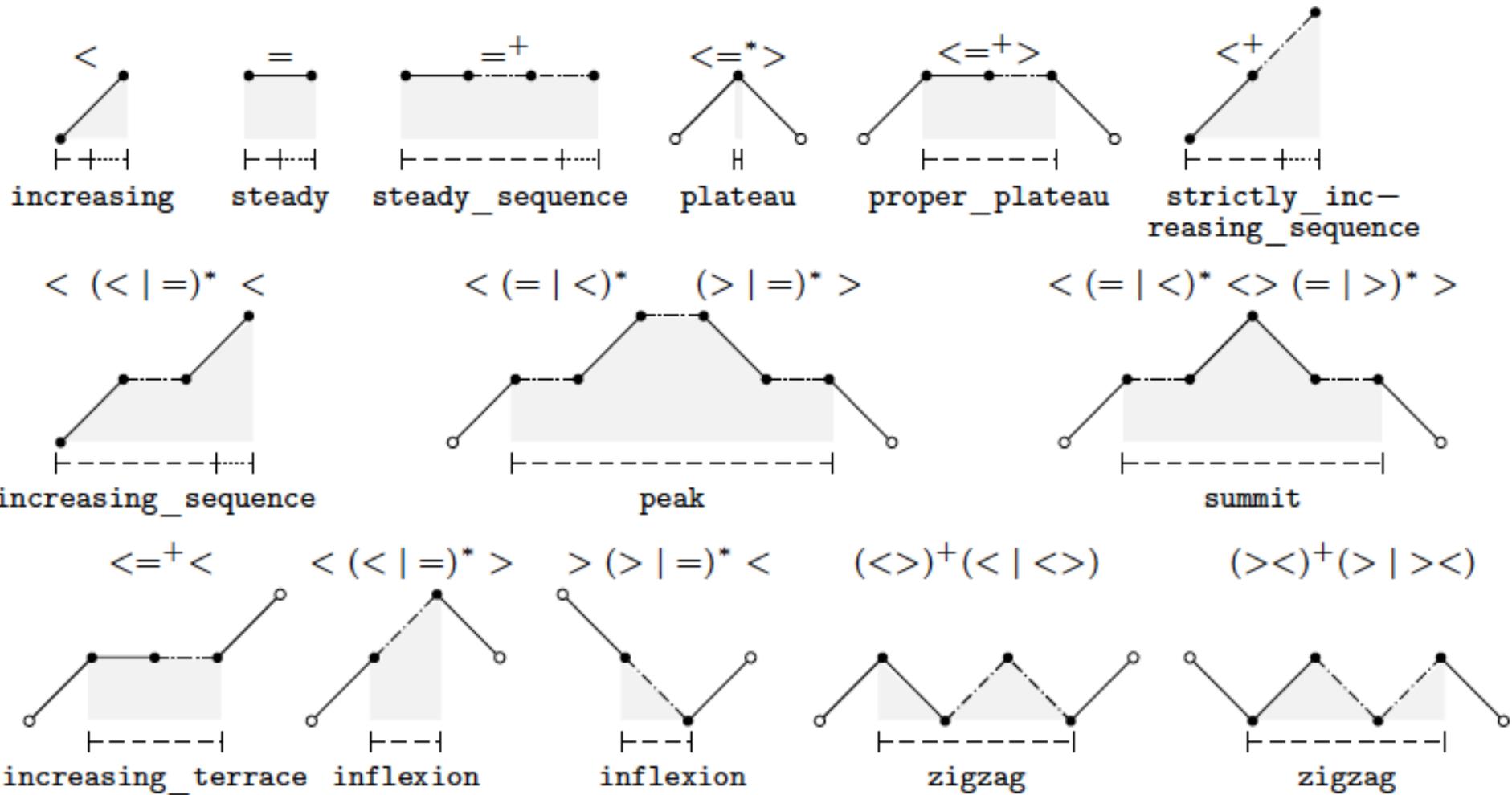
# Patterns

- **Identify a pattern** we are looking for  $r$
- **Extract subpart** for which computes feature  $a, b$

pattern	regular expression $r$	before $b$	after $a$
increasing	$<$	0	0
increasing_sequence	$< (<   =)^* <   <$	0	0
increasing_terrace	$< =^+ <$	1	1
summit	$(<   (< (=   <)^* <)) (>   (> (=   >)^* >))$	1	1
plateau	$< =^* >$	1	1
proper_plateau	$< =^+ >$	1	1
strictly_increasing_sequence	$<^+$	0	0
peak	$< (=   <)^* (>   =)^* >$	1	1
inflexion	$< (<   =)^* >   > (>   =)^* <$	1	1
steady	$=$	0	0
steady_sequence	$=^+$	0	0
zigzag	$(<>)^+ (<   <>)   (><)^+ (>   ><)$	1	1

Find **maximal** words matching regular expression  $r$

# Patterns (continued)



# Definition of $\{s|i|e\}$ -occurrences of an occurrence of pattern

Given

an input sequence  $x_0, x_1, \dots, x_{n-1}$ ,

its signature sequence  $s_0, s_1, \dots, s_{n-2}$ ,

a pattern  $(r, a, b)$ ,

a non-empty signature subsequence  $s_i, s_{i+1}, \dots, s_j$

forming a maximum word matching  $r$

the  $s$ -occurrence  $(i..j)$  is the index sequence  $i, \dots, j$ ,

the  $i$ -occurrence  $[(i+b)..j]$  is the index sequence  $i+b, \dots, j$ ,

the  $e$ -occurrence  $[(i+b)..(j+1-a)]$  is the index sequence

$i+b, \dots, j+1-a$ .

# Definition of $\{s|i|e\}$ -occurrences of an occurrence of pattern

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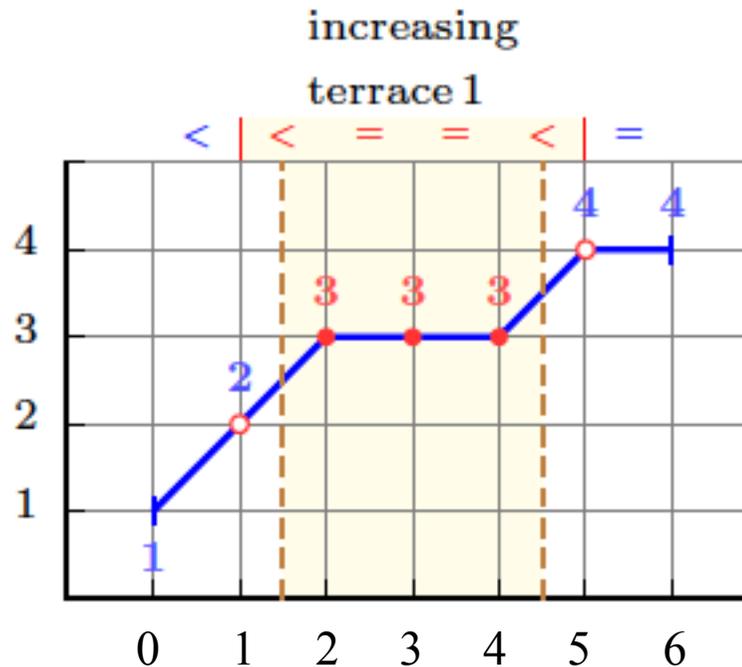
$s$ -occurrences : maximal signature sequence matching  $r$

$i$ -occurrences : do not overlap (footprint of the pattern)

$e$ -occurrences : used to compute the feature value

# Example of $\{s|i|e\}$ -occurrences for the `increasing_terrace` pattern

pattern	regular expression $r$	before $b$	after $a$
increasing terrace	$<=^+ <$	1	1



## Indices of

s-occurrences : **1..4** ( $<==<$ )

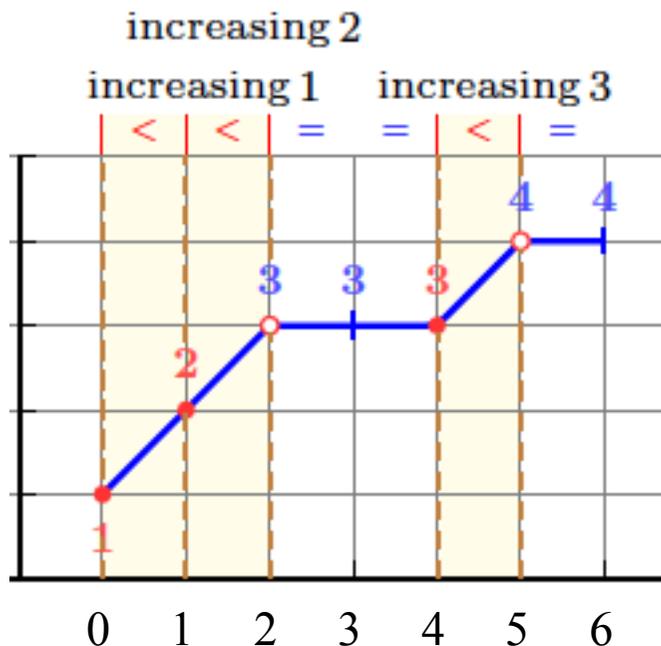
$i$ -occurrences : **[2..4]** (333)

e-occurrences : **[[2..4]]** (333)

$a=1=b$  excludes first and last input values **2** and **4**

# Example of $\{s|i|e\}$ -occurrences for the **increasing** pattern

pattern	regular expression $r$	before $b$	after $a$
increasing	$<$	0	0



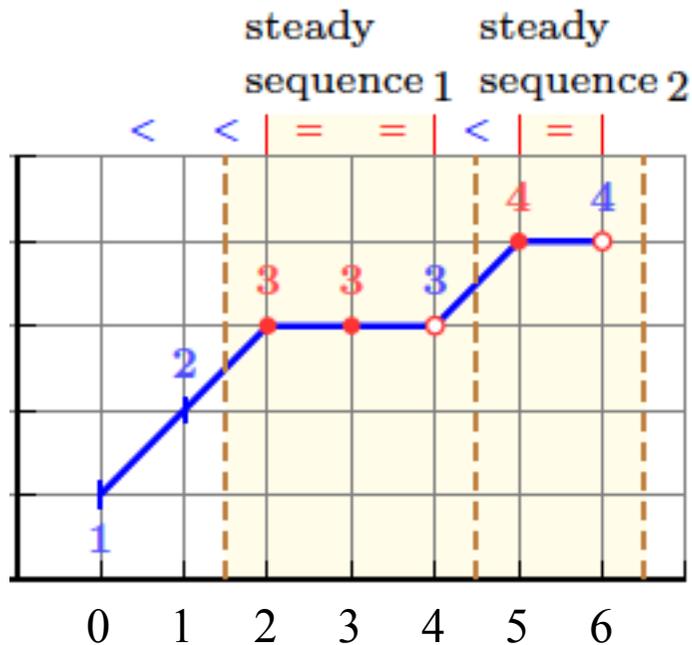
## Indices of

s-occurrences : **0..0**    **1..1**    **4..4**  
 i-occurrences : **[0..0]**    **[1..1]**    **[4..4]**  
 e-occurrences : **[[0..1]]**    **[[1..2]]**    **[[4..5]]**

since  $b=0$  s-occurrences and  $i$ -occurrences match  
 since  $a=0=b$  the 1 and 2 input values are part of e-occurrences

# Example of $\{s|i|e\}$ -occurrences for the **steady\_sequence** pattern

pattern	regular expression $r$	before $b$	after $a$
steady_sequence	$\cdot =^+ \cdot$	0	0



## Indices of

s-occurrences : **2..3**      **5..5**  
*i*-occurrences : **[2..3]**      **[5..5]**  
 e-occurrences : **[[2..4]]**      **[[5..6]]**

since  $b=0$  s-occurrences and *i*-occurrences match  
 since  $a=0=b$  the 1 and 2 input values are part of e-occurrences

# Features (*computed from e-occurrences*)

- `one` : value 1
- `width` : number of positions of the e-occurrence
- `surf` : sum of the values of the e-occurrence
- `max` : maximum value of the e-occurrence
- `min` : minimum value of the e-occurrence
- `range` : range of the e-occurrence: `max-min`

# Aggregators (*computed from sequence of features*)

- `max` : largest value of a sequence of features
- `min` : smallest value of a sequence of features
- `sum` : sum of the features of a sequence of feature

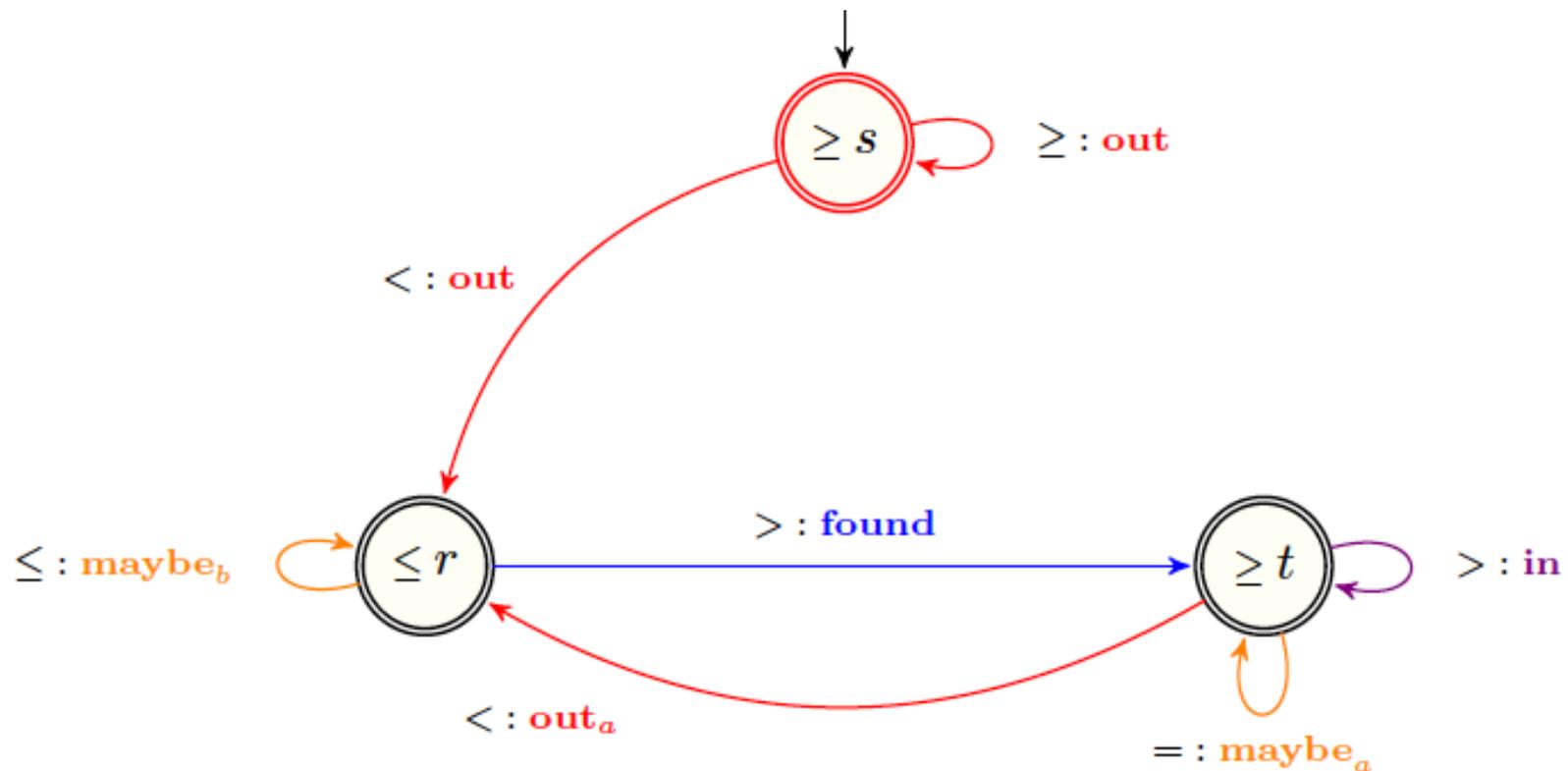
# Device for recognizing *i*-occurrences of a pattern: a seed transducer

- Define a pattern by a **transducer** (*reading/writing regular language*)
- **Input:** signature sequence
- **Output:** word of a semantic alphabet with letters:
  - **out** we are **outside** the pattern
  - **maybe** we are **possibly in** the pattern (*must be confirmed later on*)
  - **found** **first place** we know **we are in** the pattern
  - **in** we are **still** in the pattern

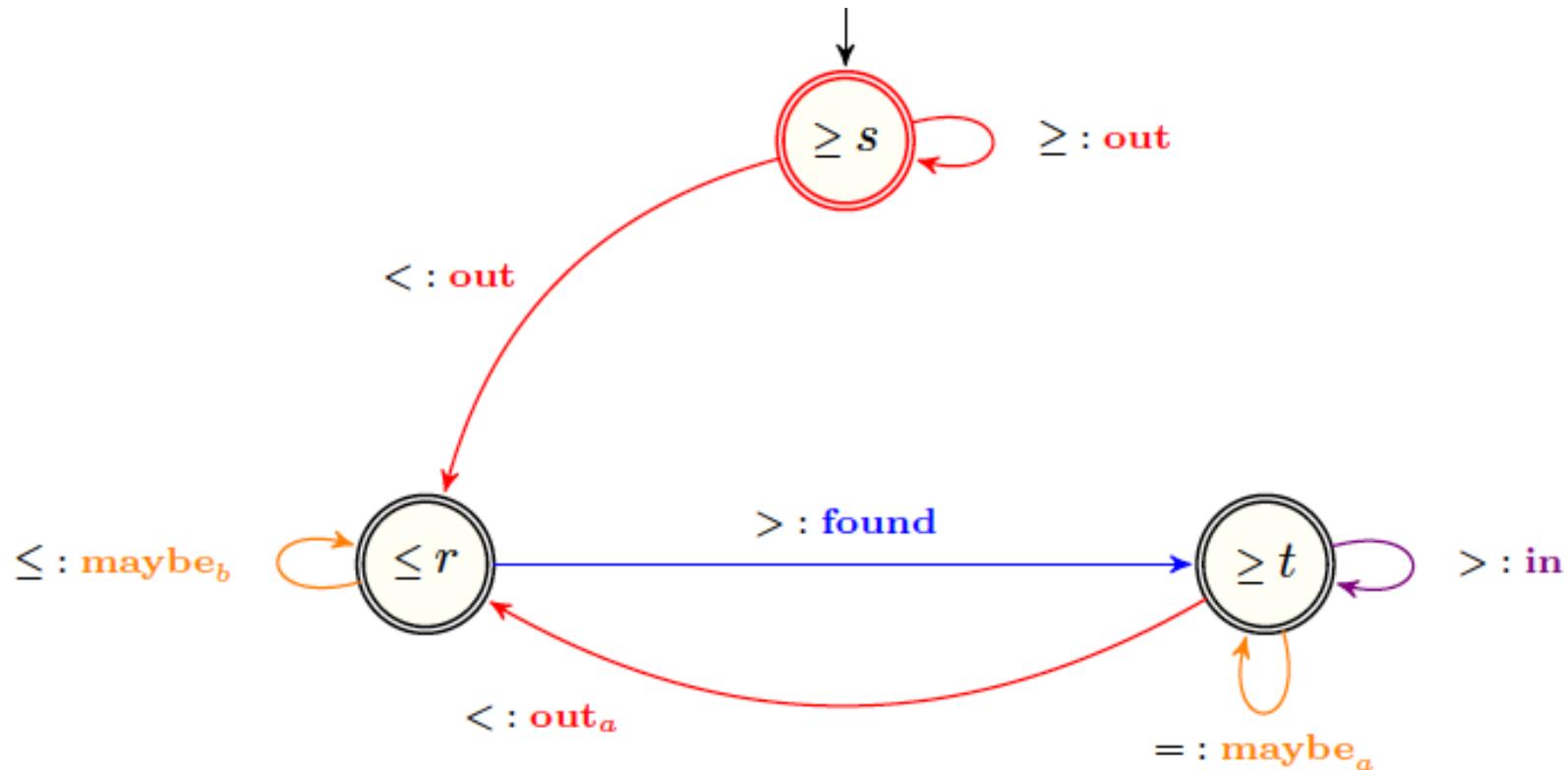
**everything** *will be synthesized from the seed transducer*

# Example: transducer for the **peak** pattern

pattern	regular expression $r$	before $b$	after $a$
peak	$\langle (=   \langle)^* (\rangle   =)^* \rangle$	1	1



# Example: transducer for the **peak** pattern



	o	o	o	o	m <sub>b</sub>	m <sub>b</sub>	m <sub>b</sub>	f	in	m <sub>a</sub>	o <sub>a</sub>	m <sub>b</sub>	f	output: semantic string				
	=	>	=	<	<	=	<	>	>	=	<	=	=	=	=	=	>	input: signature string
states:	s	s	s	s	r	r	r	r	t	t	t	r	r	r	r	r	t	input: integer sequence
	4	4	2	2	3	5	5	6	3	1	1	2	2	2	2	2	2	

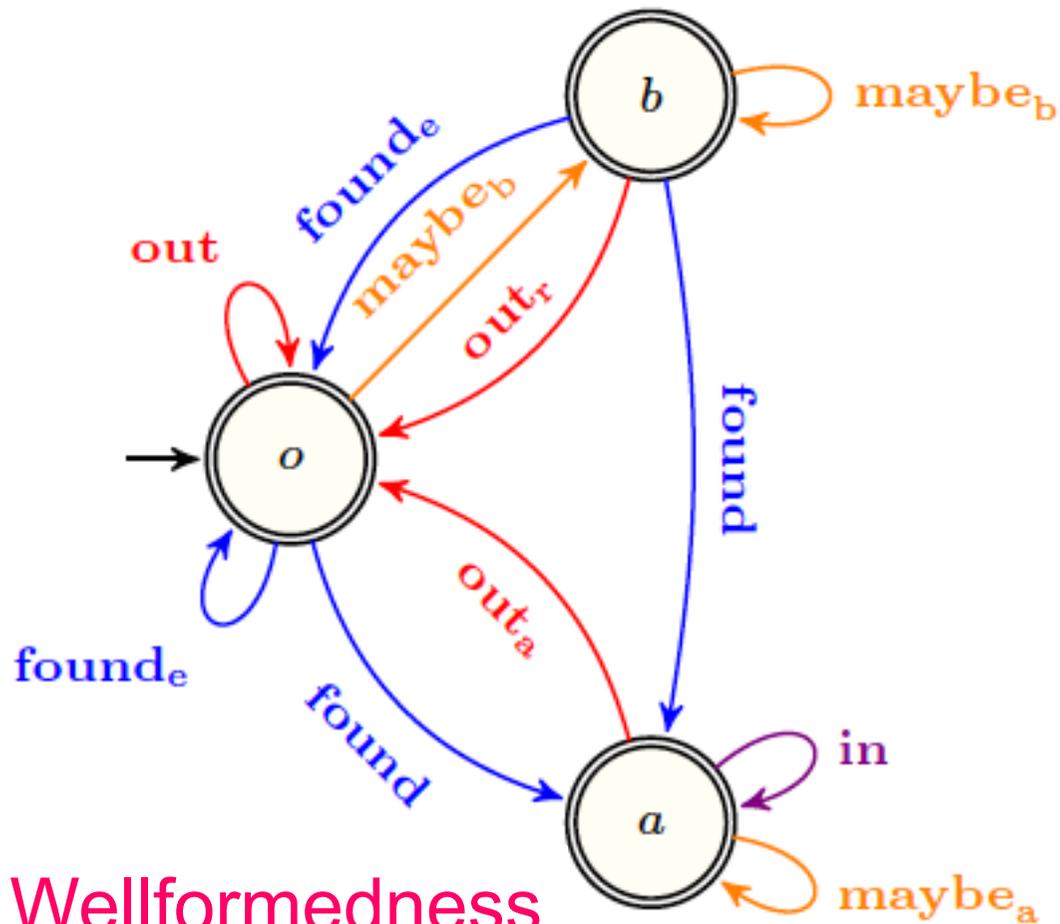
# Well-formed seed transducer (*language of the output*)

state semantics

$o$  : outside or after the end of a pattern

$b$  : potentially inside (before a **found**/**found<sub>e</sub>**)

$a$  : potentially inside (after a **found**)



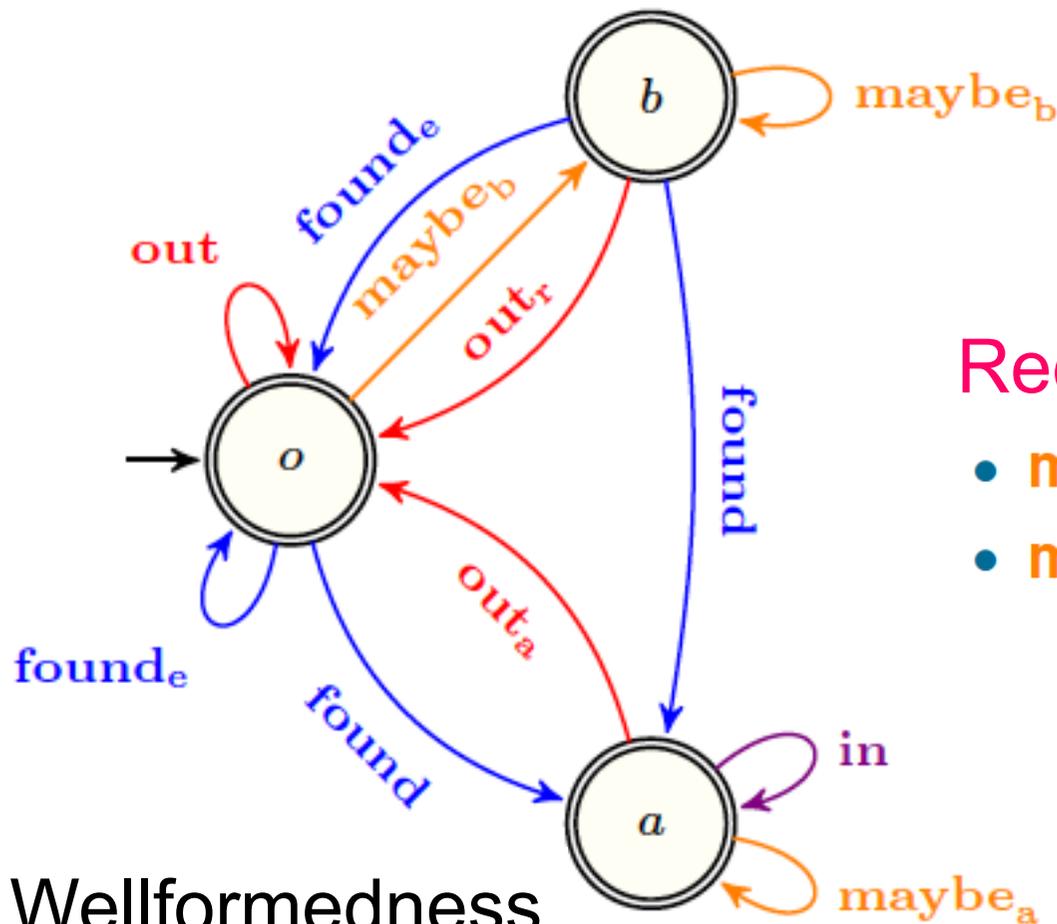
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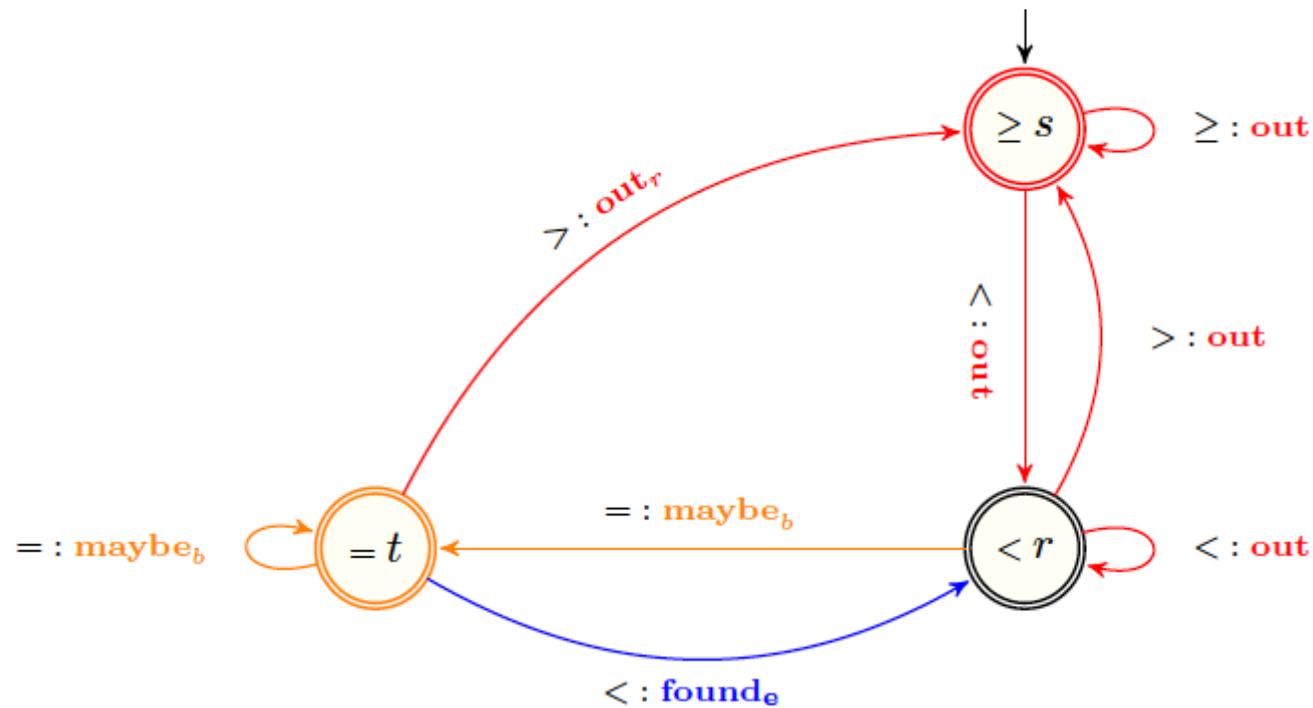
Recognizing pattern

- $\text{maybe}_b^* \text{found}_e$
- $\text{maybe}_b^* \text{found} \{ \text{maybe}_a^* \text{in}^+ \}^*$

Wellformedness

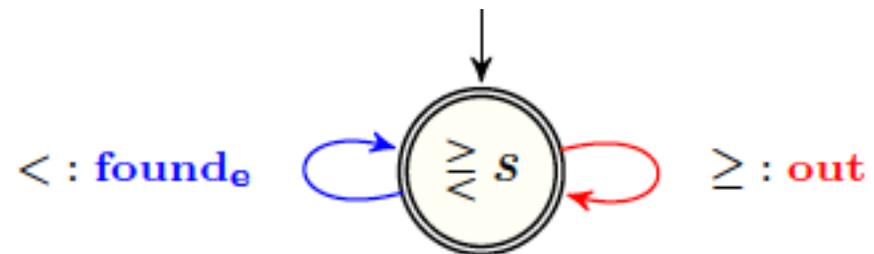
# Transducer for `increasing_terrace`

pattern	regular expression $r$	before $b$	after $a$
increasing terrace	$<=^+ <$	1	1



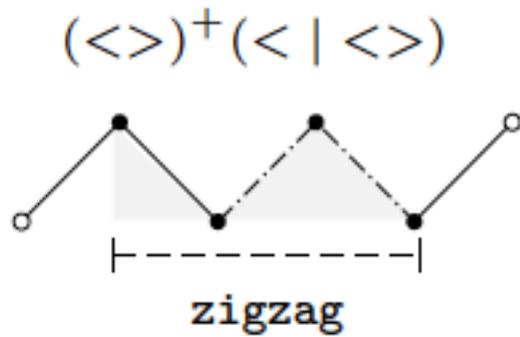
# Transducer for increasing

pattern	regular expression $r$	before $b$	after $a$
increasing	$<$	0	0

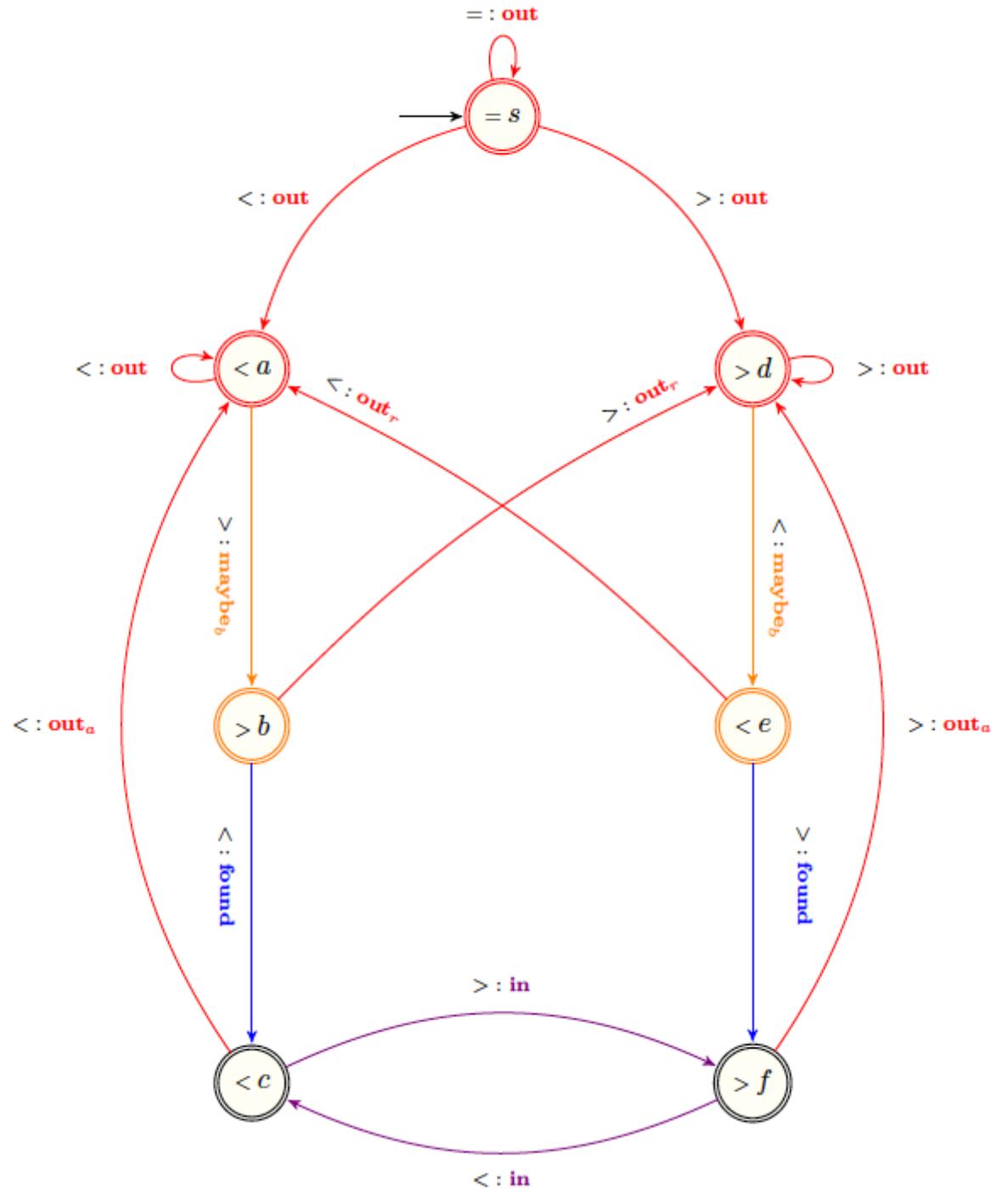
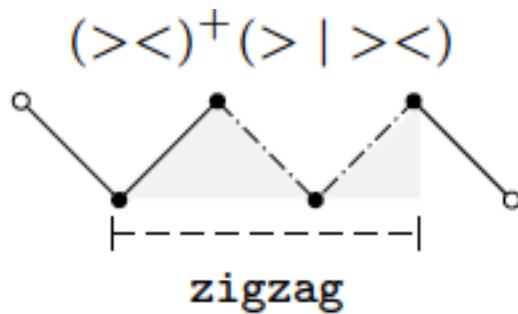


# Transducer for zigzag

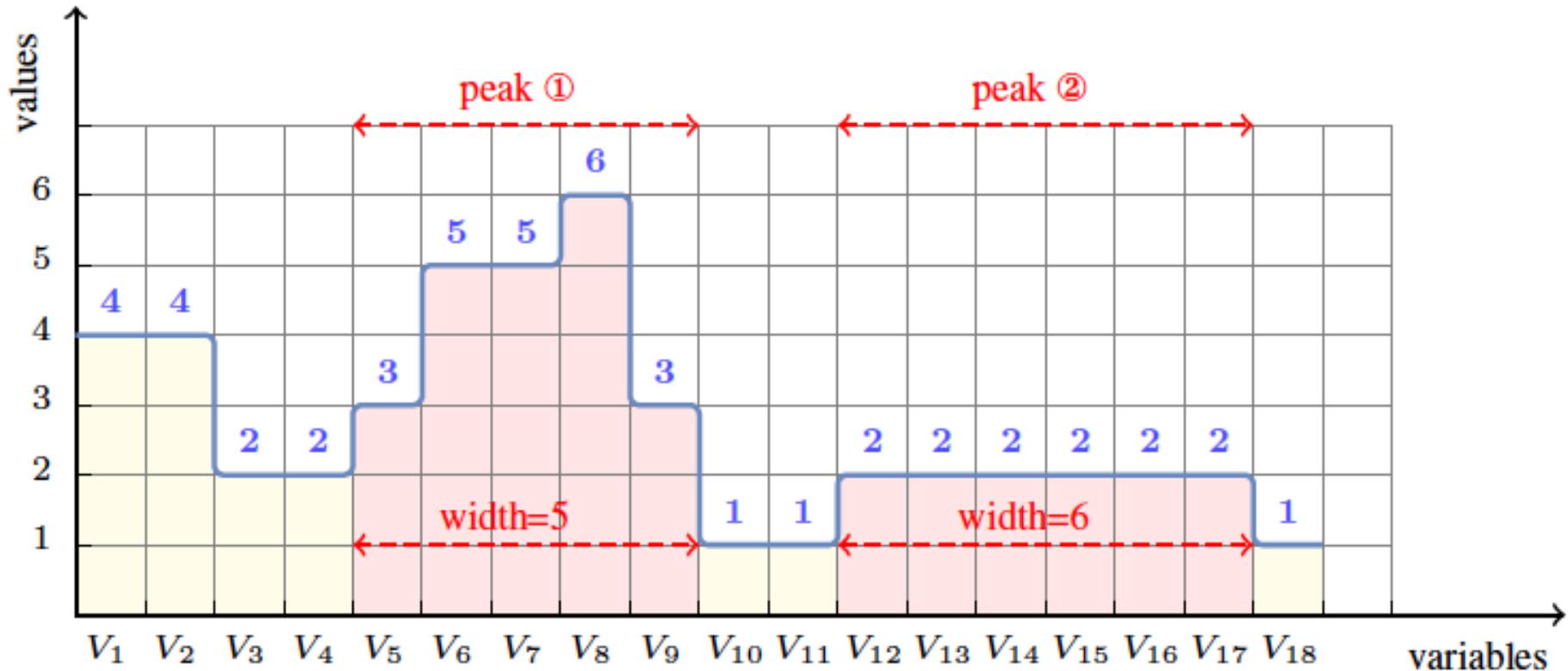
$a=1=b$



or



# Footprint constraint: identifying $i$ -occurrences of a pattern

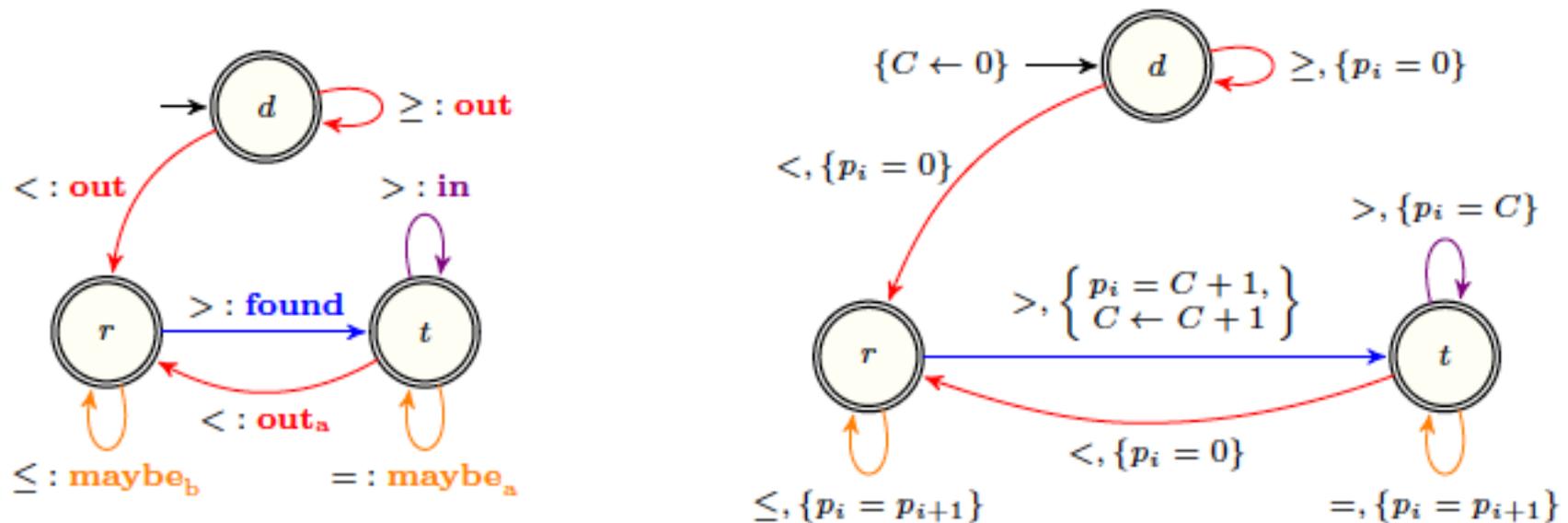


footprint(peak, [4,4,2,2,3,5,5,6,3,1,1,2,2,2,2,2,2,1],  
[0,0,0,0,1,1,1,1,1,0,0,2,2,2,2,2,2,0] )

# Decoration table for the footprint constraint (*generating counters updates*)

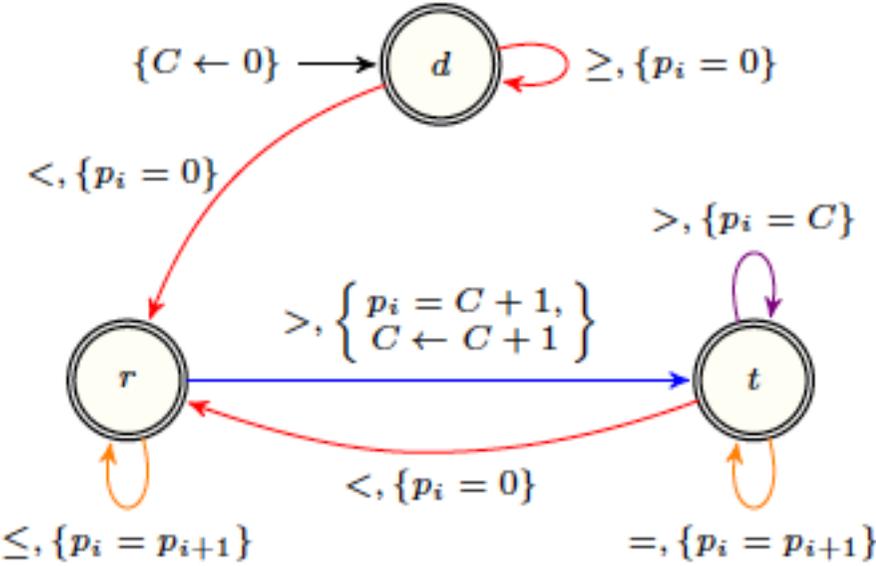
<b>initialisation</b>	$C \leftarrow 0$	
<b>return</b>	$p_n = 0$	
<b>semantic letters</b>	<b>annotations</b>	
	<i>guard</i>	<i>update of C</i>
<b>out</b>	$p_i = 0$	
<b>out<sub>r</sub></b>	$p_i = 0$	
<b>out<sub>a</sub></b>	$p_i = 0$	
<b>maybe<sub>b</sub></b>	$p_i = p_{i+1}$	
<b>maybe<sub>a</sub></b>	$p_i = p_{i+1}$	
<b>found<sub>e</sub></b>	$p_i = C + 1 \quad C \leftarrow C + 1$	
<b>found</b>	$p_i = C + 1 \quad C \leftarrow C + 1$	
<b>in</b>	$p_i = C$	

# Example: synthesizing the footprint constraint for the **peak** pattern



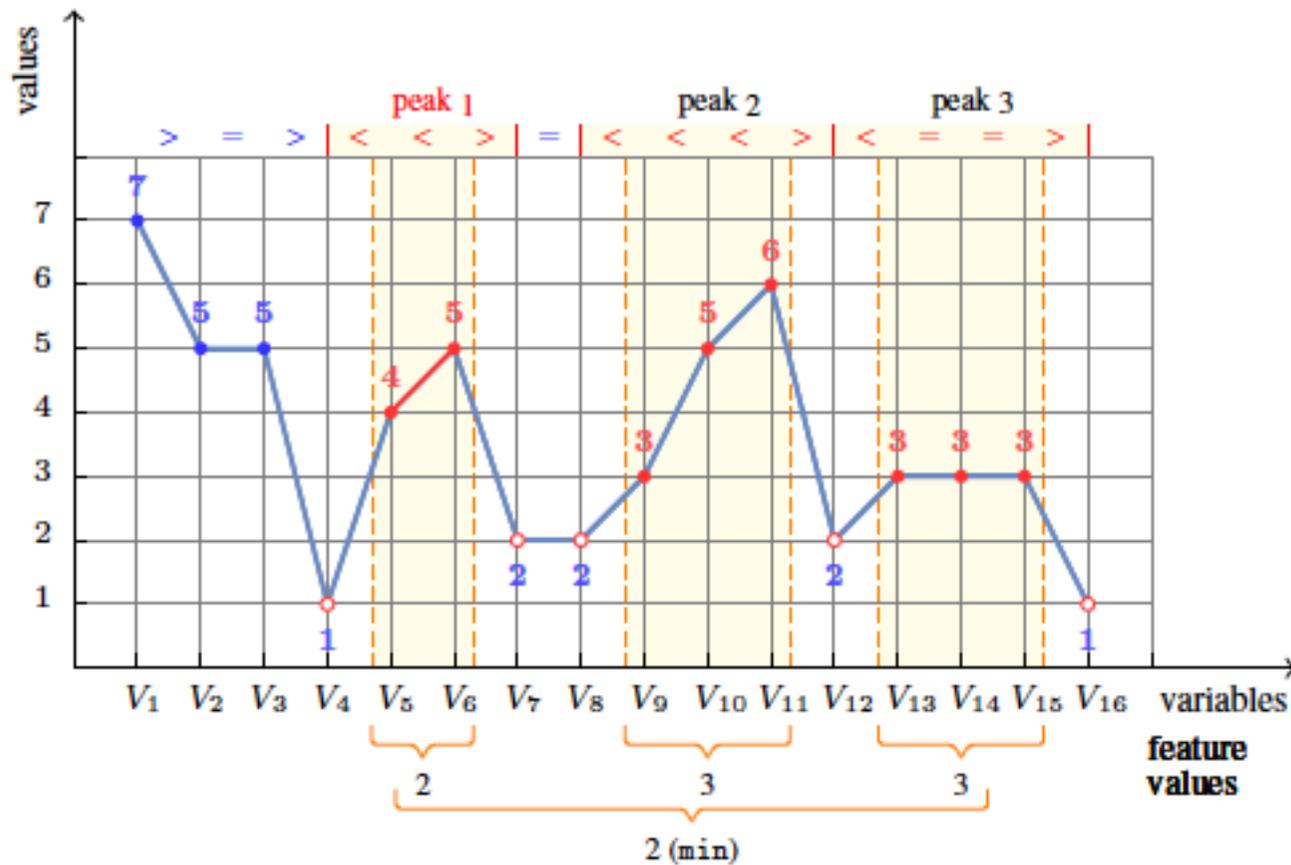
<b>out</b>	$p_i = 0$
<b>out<sub>r</sub></b>	$p_i = 0$
<b>out<sub>a</sub></b>	$p_i = 0$
<b>maybe<sub>b</sub></b>	$p_i = p_{i+1}$
<b>maybe<sub>a</sub></b>	$p_i = p_{i+1}$
<b>found<sub>e</sub></b>	$p_i = C + 1 \quad C \leftarrow C + 1$
<b>found</b>	$p_i = C + 1 \quad C \leftarrow C + 1$
<b>in</b>	$p_i = C$

# Example: executing the synthesized footprint automaton of the peak pattern



	$p_0 = 0$	$p_1 = 0$	$p_2 = 0$	$p_3 = 0$	$p_4 = p_5$	$p_5 = p_6$	$p_6 = p_7$	$p_7 = 1$	$p_8 = 1$	$p_9 = p_{10}$	$p_{10} = 0$	$p_{11} = p_{12}$	$p_{12} = p_{13}$	$p_{13} = p_{14}$	$p_{14} = p_{15}$	$p_{15} = p_{16}$	$p_{16} = 2$	$p_{17} = 0$
$C_i$	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	1	2	
$\tau_i$	o	o	o	o	$m_b$	$m_b$	$m_b$	$f$	$in$	$m_a$	$o_a$	$m_b$	$m_b$	$m_b$	$m_b$	$m_b$	$f$	
$q_i$	d	d	d	d	r	r	r	r	t	t	t	r	r	r	r	r	r	t
$s_i$	=	>	=	<	<	=	<	>	>	=	<	=	=	=	=	=	=	>
$x_i$	4	4	2	2	3	5	5	6	3	1	1	2	2	2	2	2	2	1
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17

# Feature constraints (example)



MIN\_WIDTH\_PEAK (2, [7, 5, 5, 1, 4, 5, 2, 2, 3, 5, 6, 2, 3, 3, 3, 1])

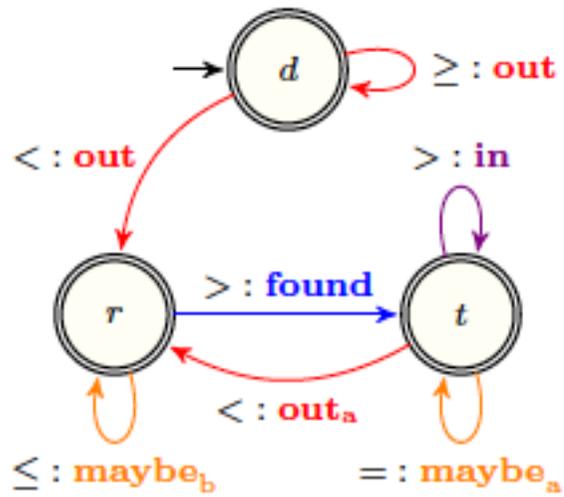
# Decoration table for the feature constraint

Feature $f$	$\text{neutral}_f$	$\text{min}_f$	$\text{max}_f$	$\phi_f$	$\delta_f$
one	1	1	1	max	0
width	0	0	$n$	+	1
surface	0	$-\infty$	$+\infty$	+	$x_t$
max	$-\infty$	$-\infty$	$+\infty$	max	$x_t$
min	$+\infty$	$-\infty$	$+\infty$	min	$x_t$
range	0	0	$+\infty$		$x_t$

Aggregator $g$	$\text{default}_{g,f}$
Max	$\text{min}_f$
Min	$\text{max}_f$
Sum	0

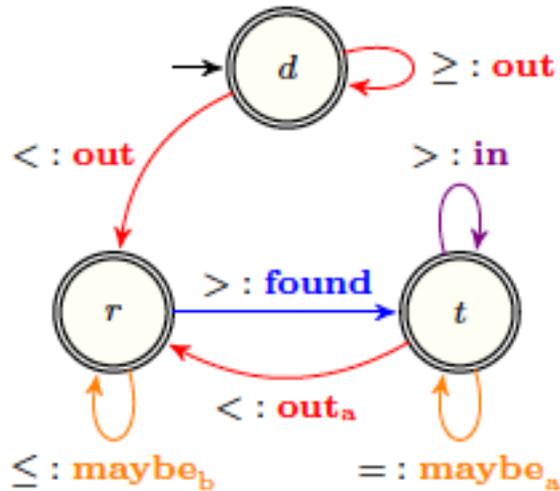
initialization		$C \leftarrow \text{default}_{af}$	$D \leftarrow \text{neutral}_f$	$R \leftarrow \text{default}_{af}$
return		$a(R, C)$		
Semantic Letter		Decoration		
	After	Update of $C$	Update of $D$	Update of $R$
out				
out <sub>r</sub>			$D \leftarrow \text{neutral}_f$	
out <sub>a</sub>		$C \leftarrow \text{default}_{af}$	$D \leftarrow \text{neutral}_f$	$R \leftarrow a(R, C)$
maybe <sub>b</sub>			$D \leftarrow \phi_f(D, \delta_f)$	
maybe <sub>a</sub>	0		$D \leftarrow \phi_f(D, \delta'_f)$	
maybe <sub>a</sub>	1		$D \leftarrow \phi_f(D, \delta_f)$	
found	0	$C \leftarrow \phi_f(\phi_f(D, \delta_f), \delta'_f)$	$D \leftarrow \text{neutral}_f$	
found	1	$C \leftarrow \phi_f(D, \delta_f)$	$D \leftarrow \text{neutral}_f$	
in	0	$C \leftarrow \phi_f(C, \phi_f(D, \delta'_f))$	$D \leftarrow \text{neutral}_f$	
in	1	$C \leftarrow \phi_f(C, \phi_f(D, \delta_f))$	$D \leftarrow \text{neutral}_f$	
found <sub>e</sub>	0		$D \leftarrow \text{neutral}_f$	$R \leftarrow a(R, \phi_f(\phi_f(D, \delta_f), \delta'_f))$
found <sub>e</sub>	1		$D \leftarrow \text{neutral}_f$	$R \leftarrow a(R, \phi_f(D, \delta_f))$

# Example: synthesizing the automaton for min\_width\_peak

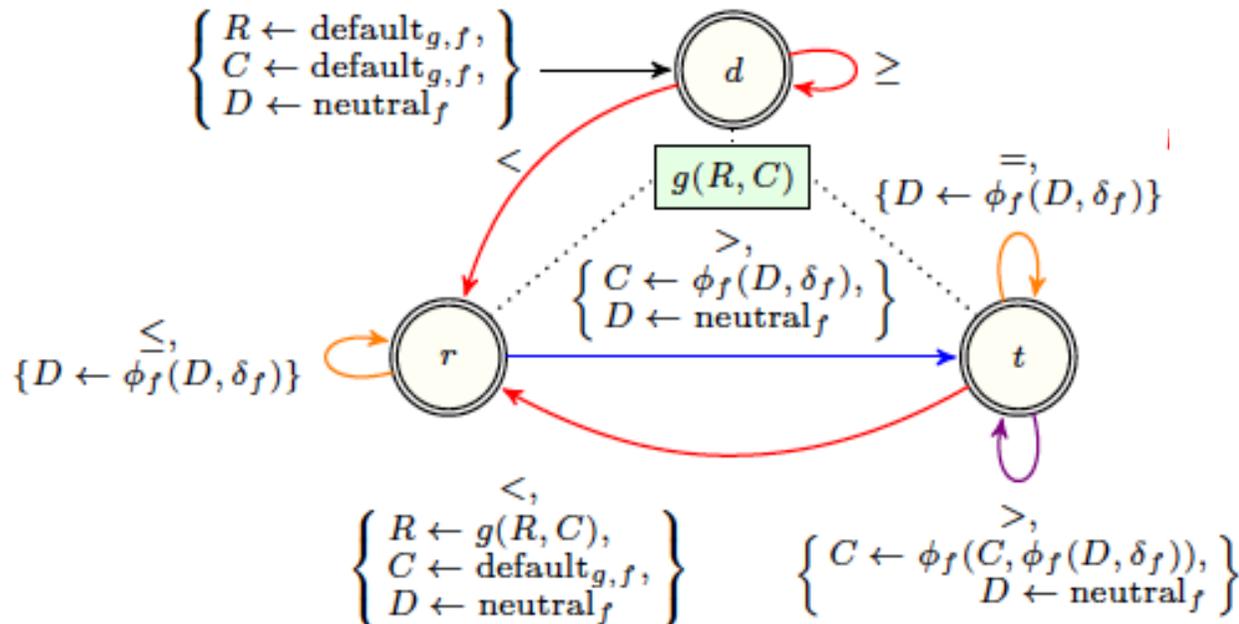


seed transducer

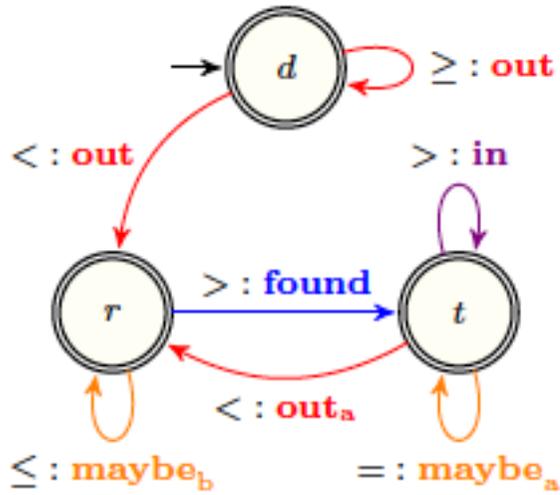
# Example: synthesizing the automaton for min\_width\_peak



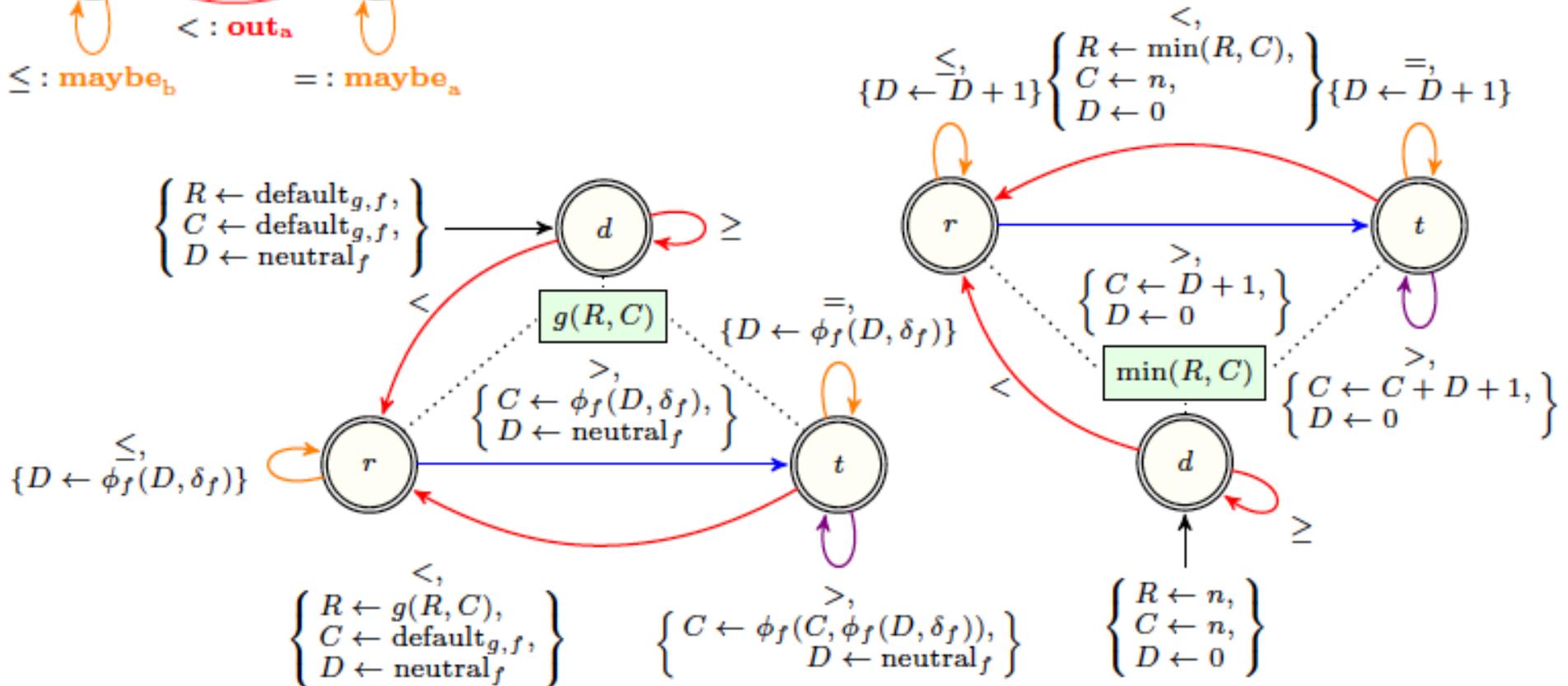
parametrized constraint  
(feature f, aggregator g)



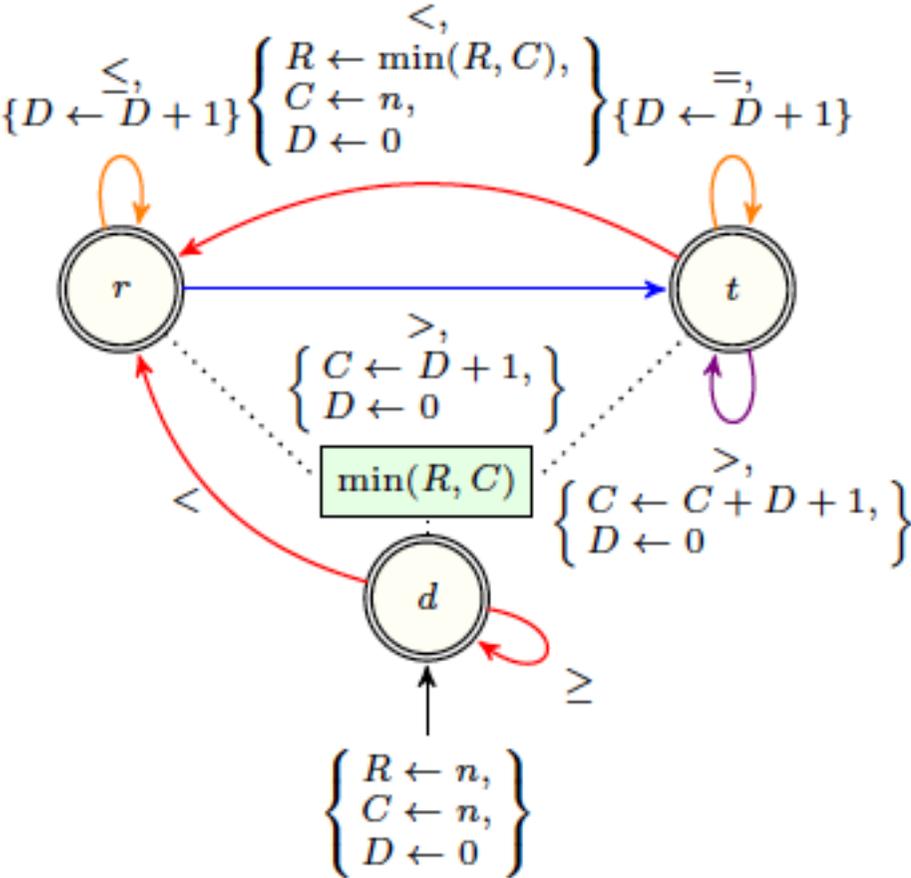
# Example: synthesizing the automaton for min\_width\_peak



specific constraint  
(min\_width\_peak)



# Running min\_width\_peak



aggregate $R$	18	18	18	18	18	18	18	18	18	18	5	5	5	5	5	5	5	min(5, 6)
current $C$	18	18	18	18	18	18	18	18	4	5	5	18	18	18	18	18	18	6
potential $D$	0	0	0	0	0	1	2	3	0	0	1	0	1	2	3	4	5	0
semantic action	$\circ$	$\circ$	$\circ$	$\circ$	$m_b$	$m_b$	$m_b$	$f$	$in$	$m_a$	$\circ_a$	$m_b$	$m_b$	$m_b$	$m_b$	$m_b$	$f$	
signatures	=	>	=	<	<	=	<	>	>	=	<	=	=	=	=	=	>	
states	$s$	$s$	$s$	$s$	$r$	$r$	$r$	$r$	$t$	$t$	$t$	$r$	$r$	$r$	$r$	$r$	$r$	$t$
input	4	4	2	2	3	5	5	6	3	1	1	2	2	2	2	2	2	1

- Background
- Synthesizing automata with accumulators from transducers
- **Parametric glue matrices**
- Simplifying automata with accumulators
- Reformulating in LP
- Deriving necessary conditions (as linear constraints)
- Bounds

# Decoration table for the feature constraint

**Problem:** since too many time-series constraints can not afford computing the glue matrix for each constraint independently

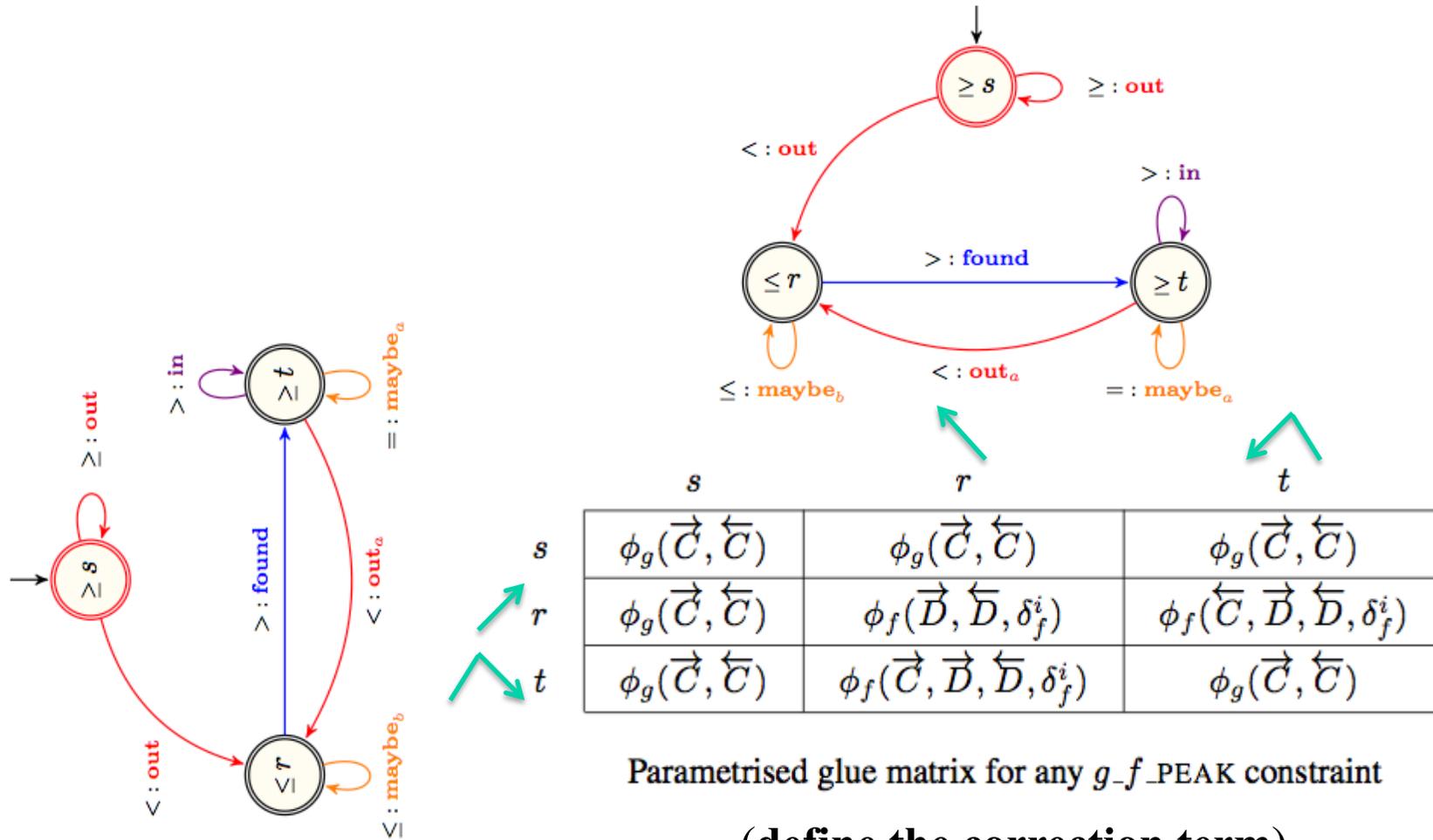
**Solution:** compute **parametrized glue matrices** at the level of the transducer (*rather than at the level of each automaton*)

# Values and functions used to parametrize generated automata (*and glue matrices*)

Feature $f$	$\text{id}_f$	$\text{min}_f$	$\text{max}_f$	$\phi_f$	$\delta_f^i$	Aggregator $g$	$\phi_g$	$\text{default}_{gf}$
one	1	1	1	1	1	Max	max	$\text{min}_f$
width	0	$n + 1$	$+\infty$	+	1	Min	min	$\text{max}_f$
surface	0	$-\infty$	$+\infty$	+	$x_i$	Sum	+	0
max	$-\infty$	$-\infty$	$+\infty$	max	$x_i$			
min	$+\infty$	$-\infty$	$+\infty$	min	$x_i$			
range	0	0	$+\infty$	n/a	$x_i$			

Table 2.1: (Left) Features: identity, minimum, and maximum values; the operators  $\phi_f$  and  $\delta_f^i$  recursively define the feature value  $v_u$  of a time series  $x_\ell, \dots, x_u$  by  $v_\ell = \phi_f(\text{id}_f, \delta_f^\ell)$  and  $v_i = \phi_f(v_{i-1}, \delta_f^i)$  for  $i > \ell$ , where  $\delta_f^i$  is the contribution of  $x_i$  to  $v_u$ ;  $n$  stands for the length of the time-series. (Right) Aggregators: operators and identity values relative a feature  $f$ .

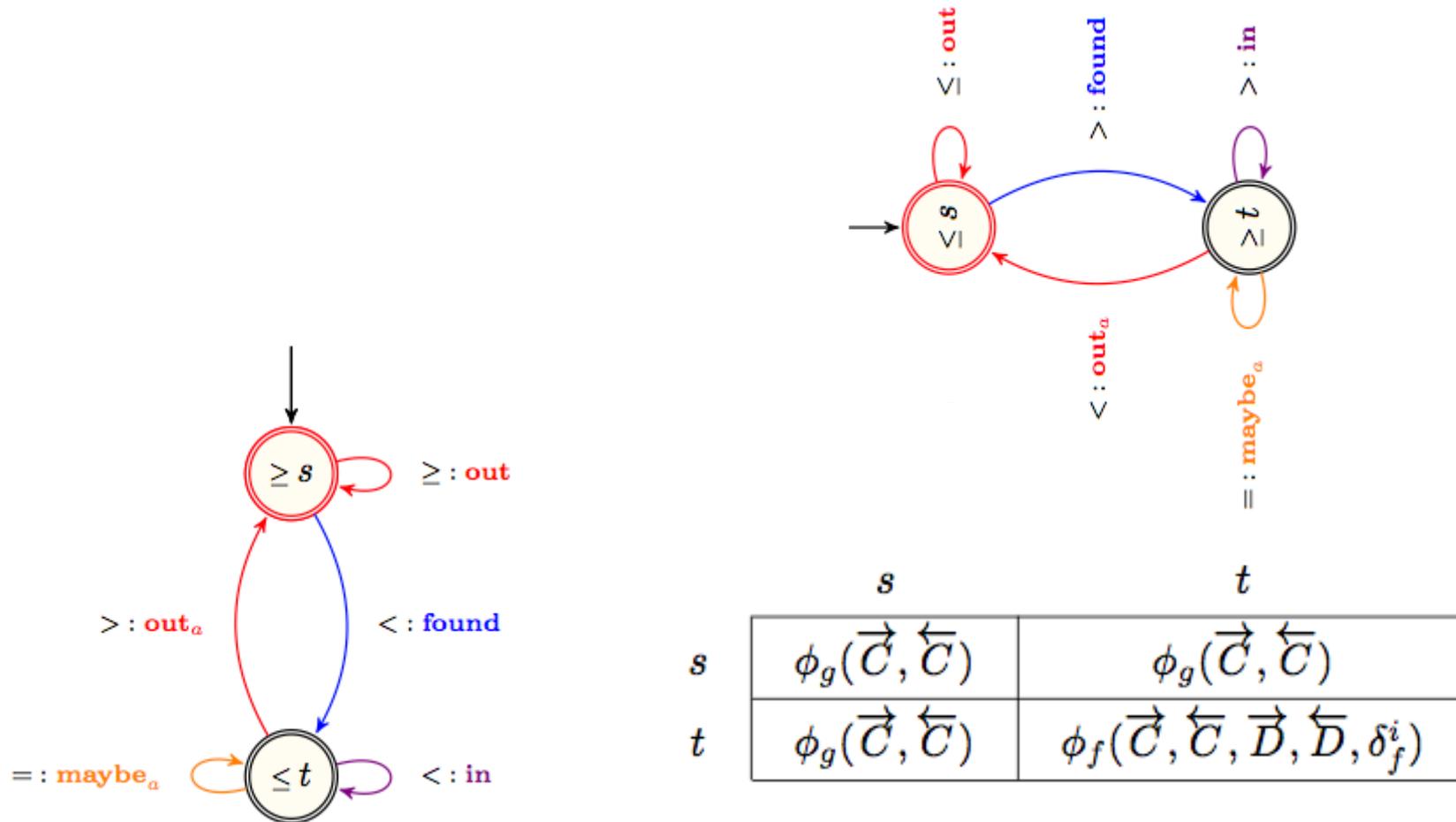
# Parametrize glue matrix for peak (is its own reverse)



Parametrised glue matrix for any  $g$ - $f$  PEAK constraint

(define the correction term)

# Parametrize glue matrix for increasing\_sequence (its reverse is decreasing\_sequence)



Parametrised glue matrix for any  $g\_f\_INCREASING\_SEQUENCE$  constraint  
(define the correction term)

- Background
- Synthesizing automata with accumulators from transducers
- Parametric glue matrices
- **Simplifying automata with accumulators**
- Reformulating in LP
- Deriving necessary conditions (as linear constraints)
- Bounds

# Automata simplifications

## Goal

- ▶ **Reduce** the number of accumulators and aggregate as **early** as possible
- ▶ Simplify the automata at the stage of their synthesis

## Three simplification types

- ▶ Simplifications coming from the properties of patterns, ex.: aggregate-once
- ▶ Simplifications coming from the properties of the feature/aggregator pairs, ex.: immediate-aggregation
- ▶ Removing the never used accumulators.

## “Aggregate-once” simplification

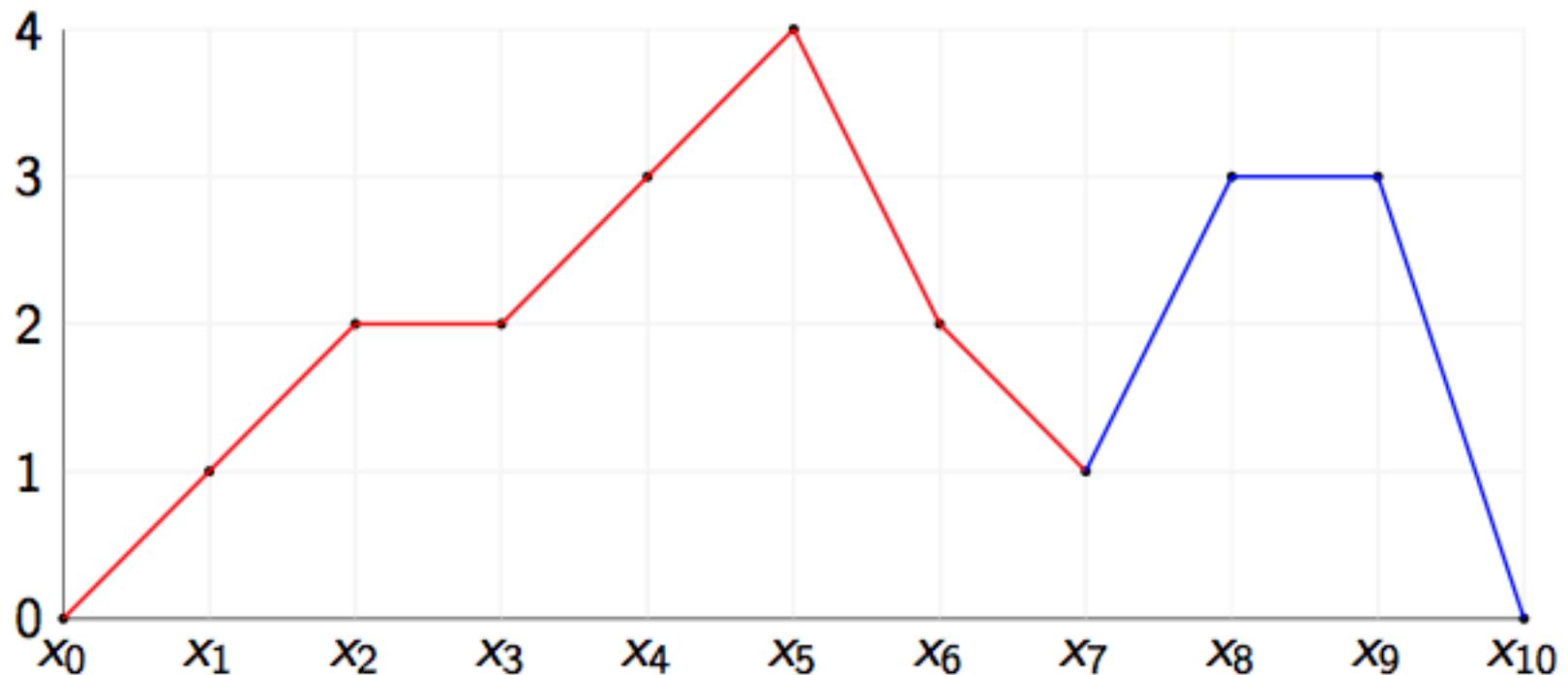
### What is the “Aggregate-once” simplification ?

It allows to compute the feature value of a current pattern occurrence only once and, possibly, earlier than the end of a pattern occurrence.

### When is the simplification applicable ?

There must exist a transition on which the value of the feature from the current pattern occurrence is known.

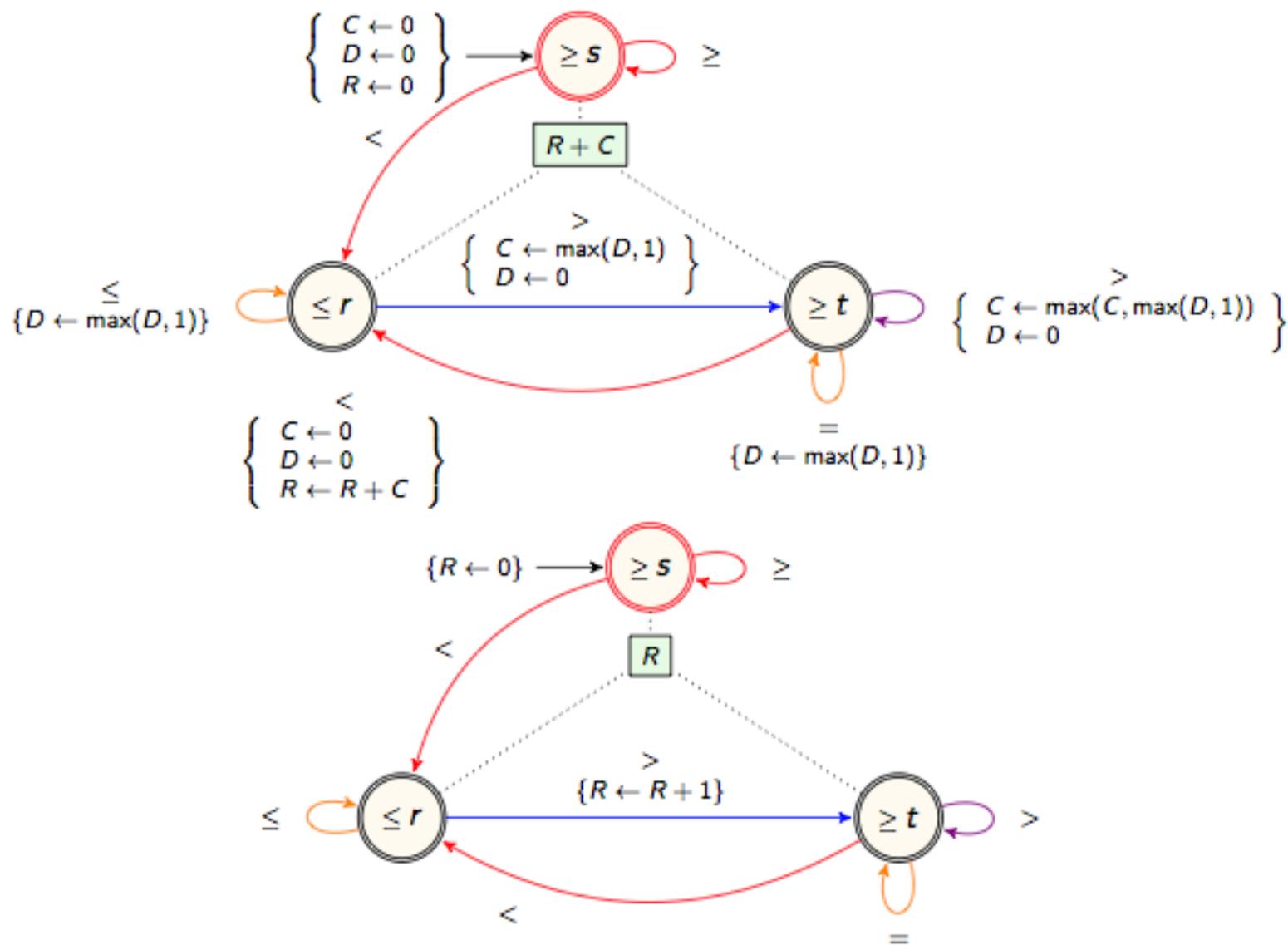
## Example: counting number of peaks



$s_0 = '<'$   $s_1 = '<'$   $s_2 = '='$   $s_3 = '<'$   $s_4 = '<'$   $s_5 = '>'$   $s_6 = '>'$   $s_7 = '<'$   $s_8 = '='$   $s_9 = '>'$

1. First peak is detected upon consuming  $s_5$
2. Second peak is detected upon consuming  $s_9$

## Two automata for nb\_peak



## Percentage of simplified constraints

Simplification	Percentage
aggregate once	28.9 %
immediate aggreg.	45.9 %
other properties	11.6 %
unchanged automata	13.6 %

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## Goal

### Goal

A way to generate a model for an automaton with linear or linearisable accumulator updates, for example containing min and max.

### Linear decomposition of automata without accumulators

Côté, M.C., Gendron, B., Rousseau, L.M.: Modeling the regular constraint with integer programming. In: CPAIOR 2007. LNCS, vol. 4510, pp. 29–43. Springer (2007)

## Signature constraint

Introduced variables:  $S_i$  over  $\Sigma$  with  $i \in [0, n - 2]$ .

What do the values of  $S_i$  mean ?

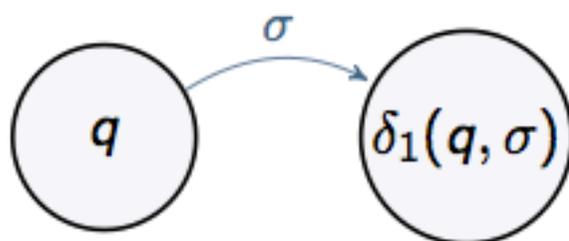
$$S_i = '>' \Leftrightarrow X_i > X_{i+1}, \forall i \in [0, n - 2]$$

$$S_i = '=' \Leftrightarrow X_i = X_{i+1}, \forall i \in [0, n - 2]$$

$$S_i = '<' \Leftrightarrow X_i < X_{i+1}, \forall i \in [0, n - 2]$$

## Transition function constraints

Introduced variables:  $Q_i$  over  $Q$  with  $i \in [0, n - 1]$ ;  $T_i$  over  $Q \times \Sigma$  with  $i \in [0, n - 2]$



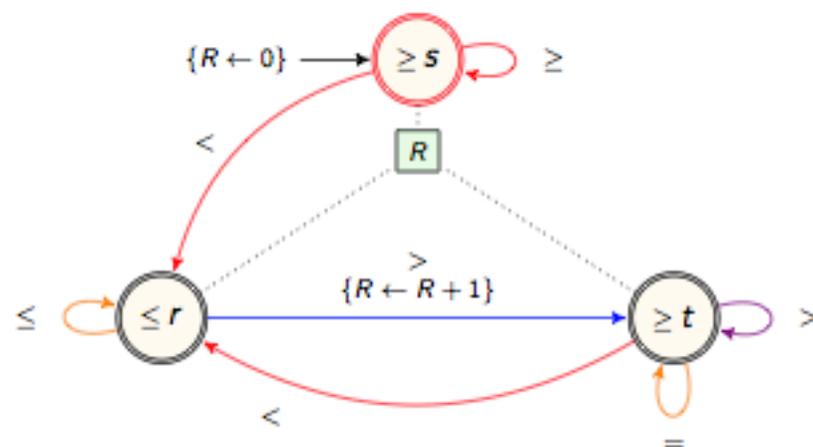
**Each transition constraint has a form:**

$$Q_i = q \wedge S_i = \sigma \Leftrightarrow Q_{i+1} = \delta_1(q, \sigma) \wedge T_i = \langle q, \sigma \rangle,$$
$$\forall i \in [0, n - 2], \forall q \in Q, \forall \sigma \in \Sigma$$

**Initial state is fixed**

$$Q_0 = q_0$$

## Accumulator updates



### Accumulator updates

$R_i$  over  $[a, b]$  with  $i$  in  $[0, n-1]$ ;  $T_i$  over  $Q \times \Sigma$  with  $i$  in  $[0, n-2]$ .

- ▶  $R_0 = 0$
- ▶  $T_i = \langle r, > \rangle \Rightarrow R_{i+1} = R_i + 1, \forall i \in [0, n-2]$
- ▶  $T_i = \langle q, \sigma \rangle \Rightarrow R_{i+1} = R_i, \forall i \in [0, n-2], \forall \langle q, \sigma \rangle \in (Q \times \Sigma) \setminus \langle r, > \rangle$
- ▶  $M = R_{n-1}$

## New variables for the linear model

### New variables

- ▶  $Q_i$  is replaced by 0-1 variables  $Q_i^q$  for all  $q$  in  $Q$ .  
 $Q_i^q = 1 \Leftrightarrow Q_i = q$
  - ▶ New constraint:  $\sum_{q \in Q} Q_i^q = 1, \forall i \in [0, \dots, n-1]$
  - ▶ The same procedure for  $T_i$  and  $S_i$  wrt their domains
  - ▶  $X_i$  and  $R_i$  remain integer variables!
- 
- ▶ Every constraint of the logical model is made linear by applying some standard techniques
  - ▶ The linear model has  $O(n)$  variables and  $O(n)$  constraints

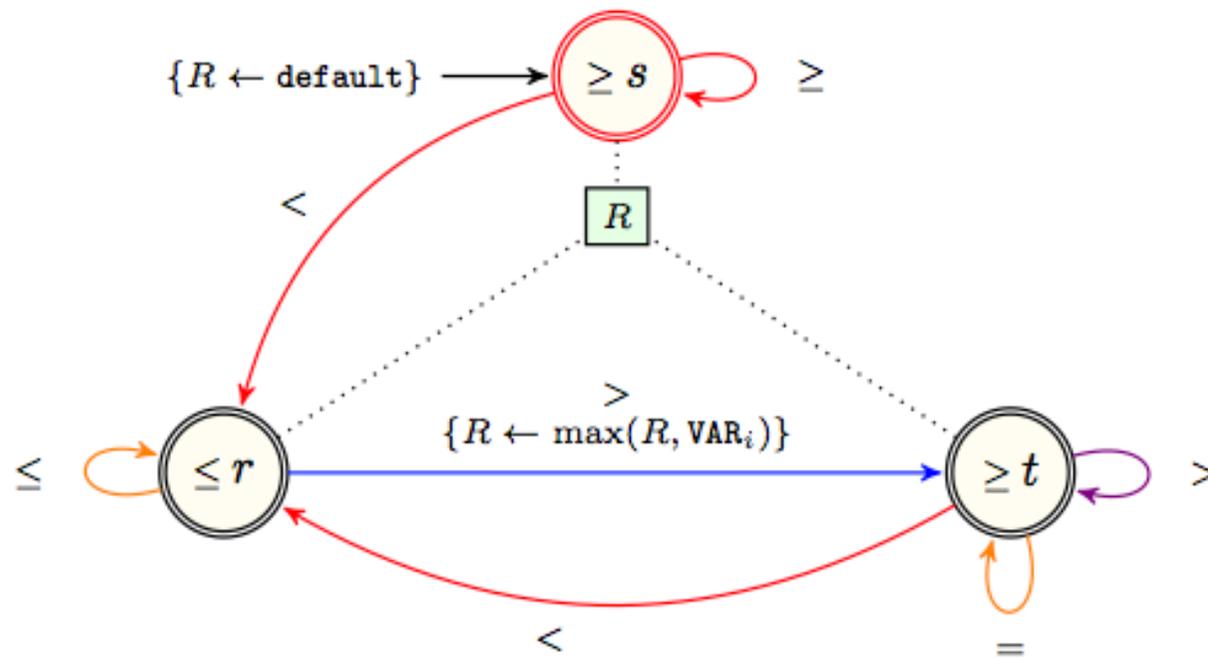
- Background
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# Generating necessary conditions (*as linear constraints*)

- Good news:
  - Can use/adapt standard LP techniques (Farkas Lemma) to generate necessary conditions (*expressed as linear constraints*)
  - Can use this even if accumulators updates are min/max operations
  - Can rank the linear constraints
  - Useful both for CP and LP
  - The invariants neither depend of the domain of the variables, nor on the size of the sequence.

Leads to a data base of cuts for time-series constraints

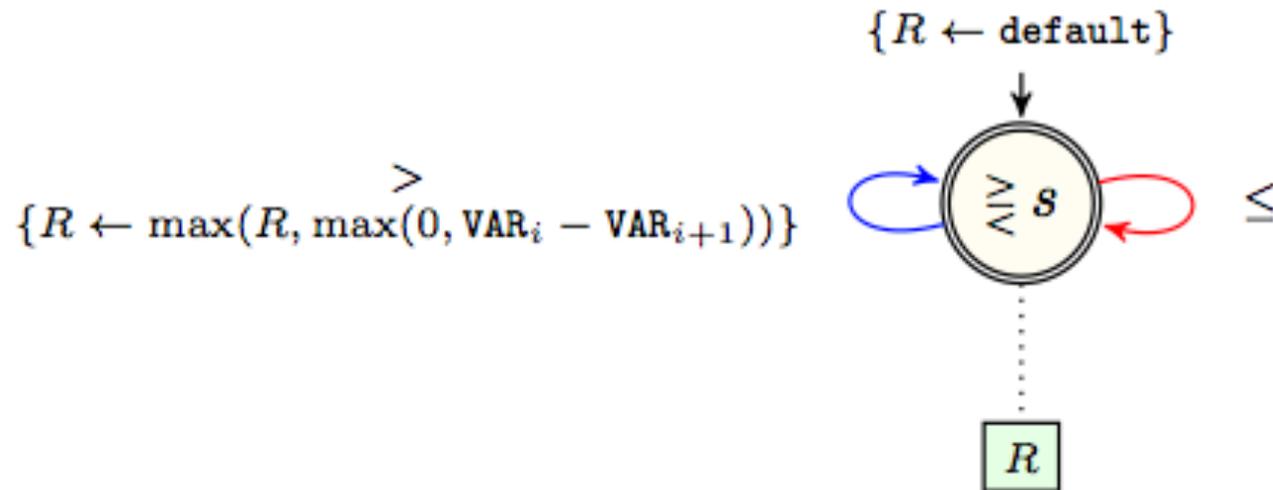
# Examples of linear constraints (max\_max\_peak)



$$R_i - R_{i-1} \geq 0$$

(increasing R since use max aggregator)

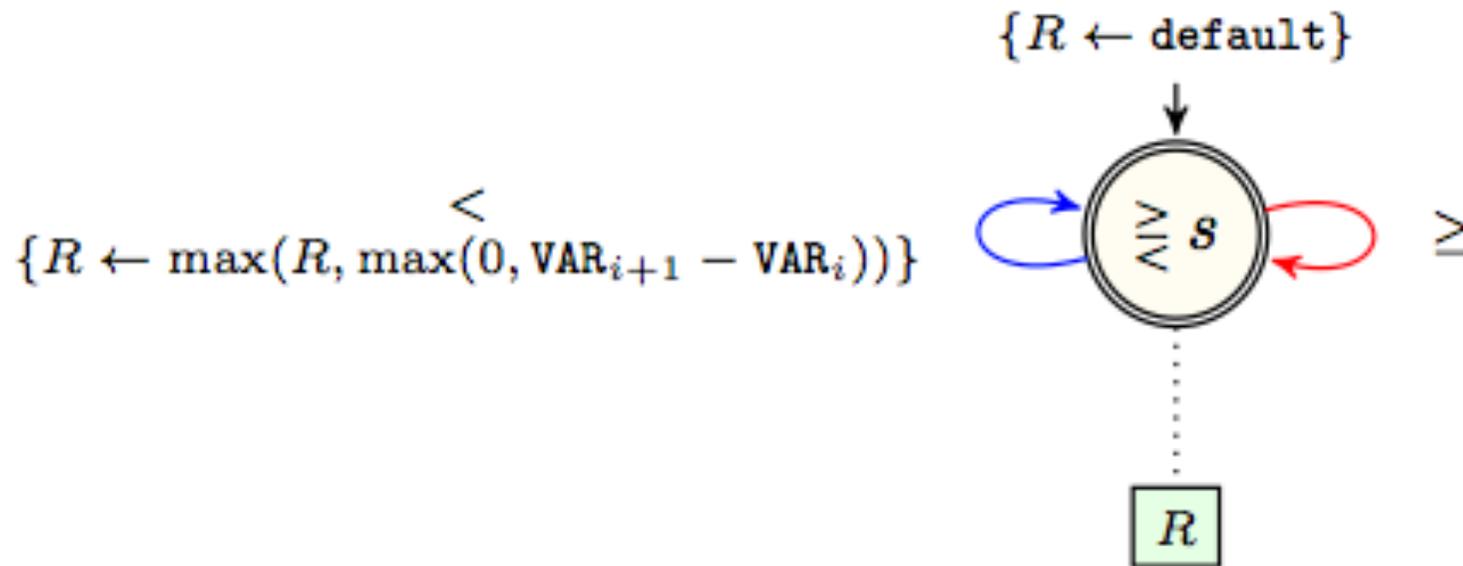
# Examples of linear constraints (max\_range\_decreasing)



$R_i - R_{i-1} \geq 0$  (increasing R since use max aggregator)

$R_i + \text{VAR}_{i-1} - \text{VAR}_{i-2} \geq 0$

# Examples of linear constraints (max\_range\_increasing)

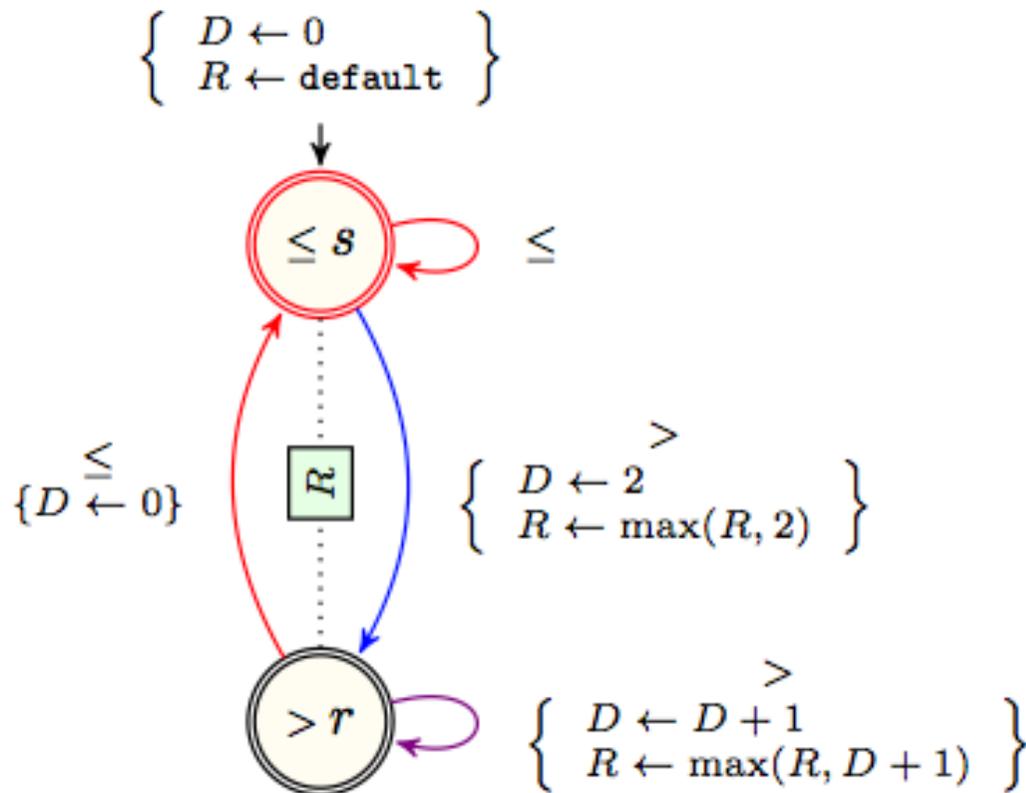


$R_i - R_{i-1} \geq 0$  (increasing R since use max aggregator)

$$R_i - \text{VAR}_{i-1} + \text{VAR}_{i-2} \geq 0$$

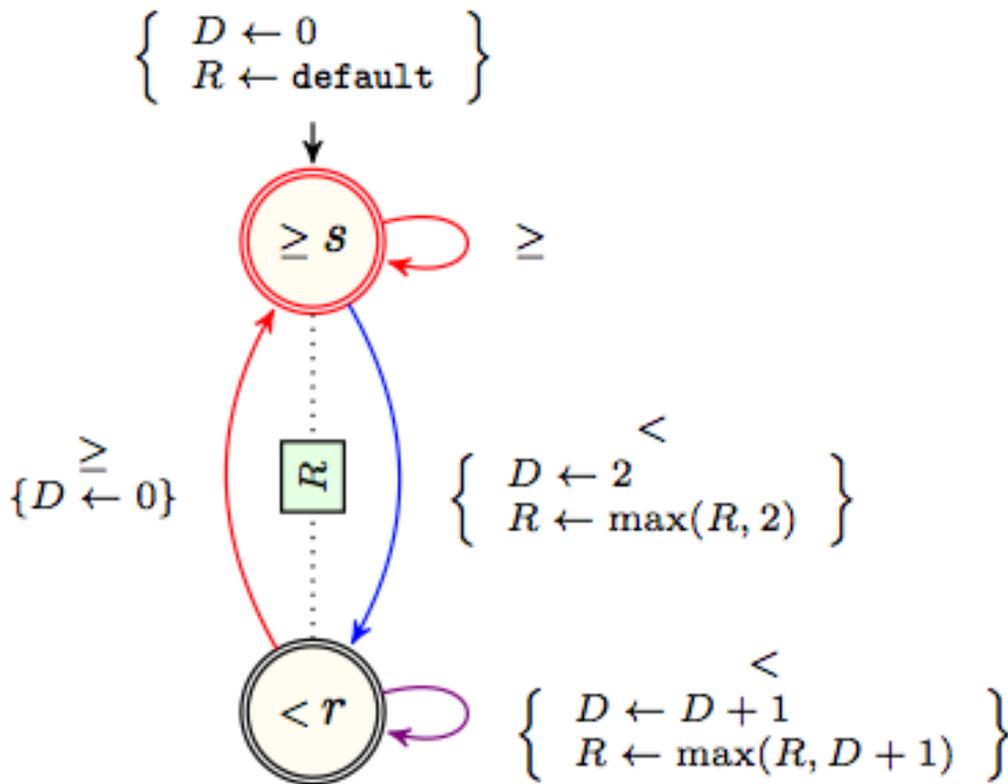
# Examples of linear constraints

(max\_width\_strictly\_decreasing\_sequence)



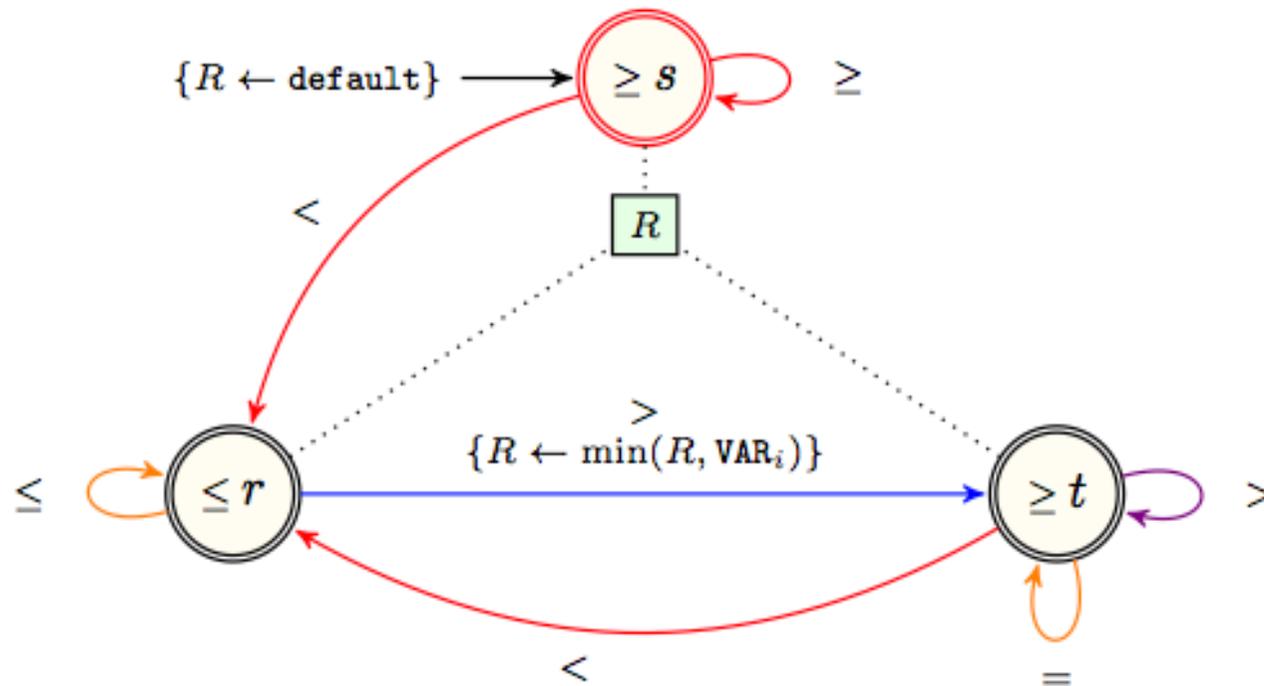
$$R_i - R_{i-1} \geq 0 \quad (\text{increasing } R \text{ since use max aggregator})$$

# Examples of linear constraints (max\_width\_strictly\_increasing\_sequence)



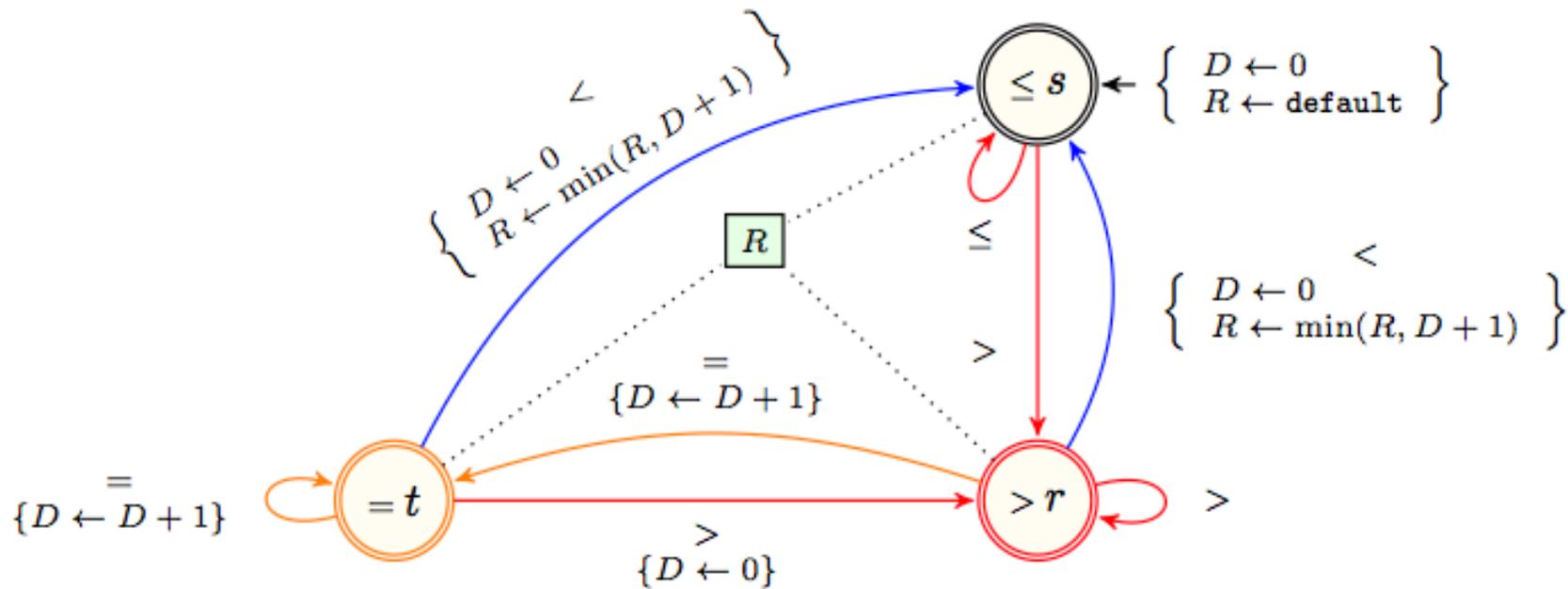
$$R_i - R_{i-1} \geq 0 \quad (\text{increasing } R \text{ since use max aggregator})$$

# Examples of linear constraints (min\_max\_peak)



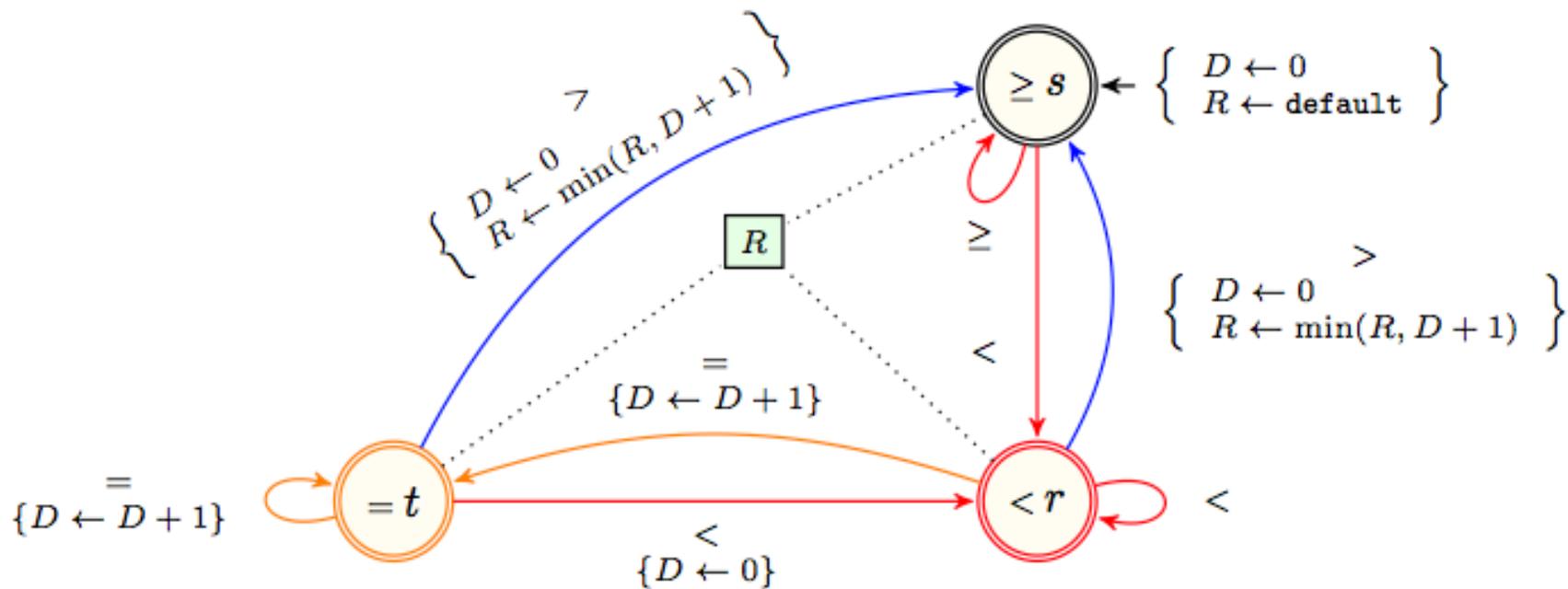
$$-R_i + R_{i-1} \geq 0 \quad (\text{decreasing } R \text{ since use min aggregator})$$

# Examples of linear constraints (min\_width\_plain)



$$-R_i + R_{i-1} \geq 0 \quad (\text{decreasing } R \text{ since use min aggregator})$$

# Examples of linear constraints (min\_width\_plateau)

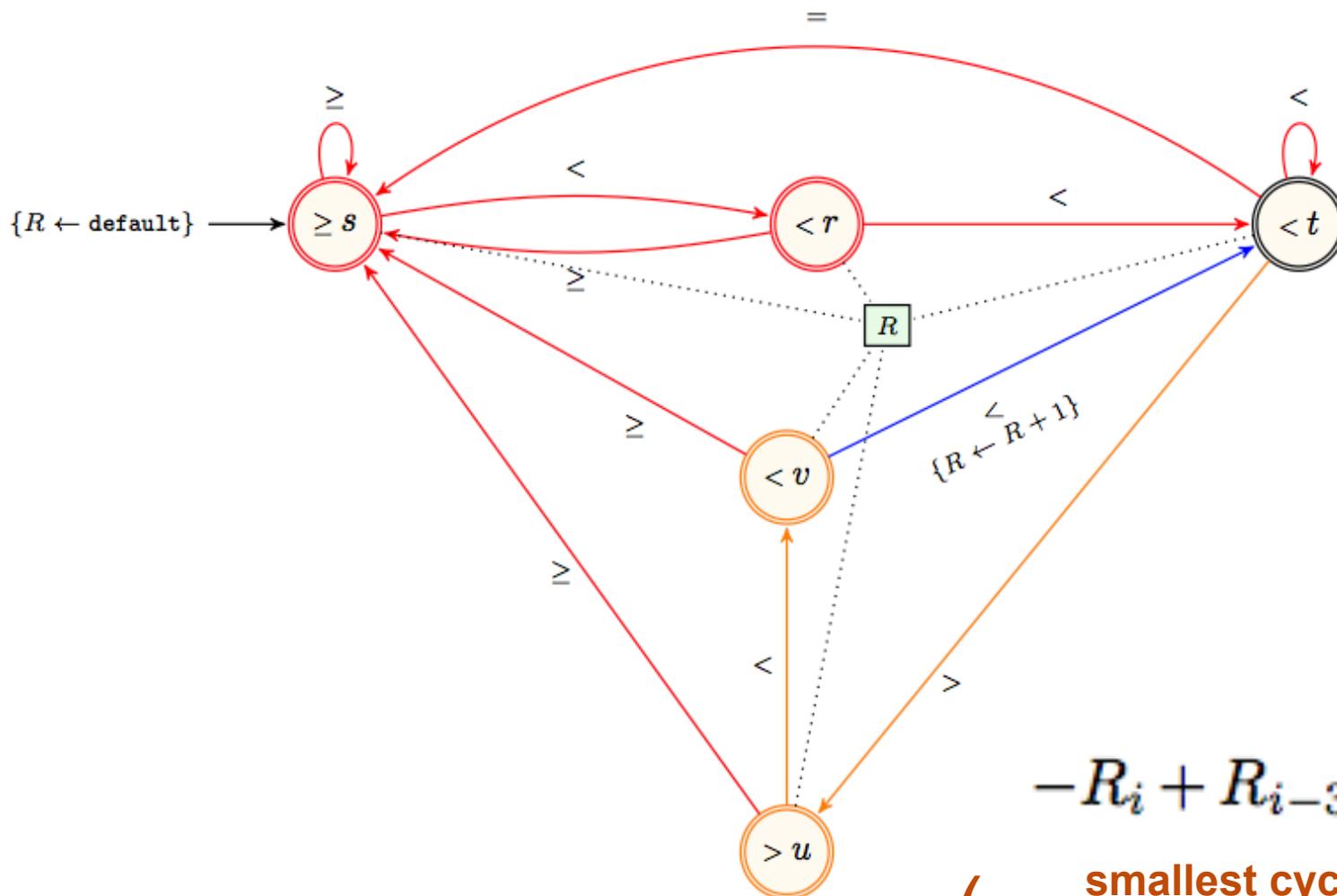


$$-R_i + R_{i-1} \geq 0 \quad (\text{decreasing } R \text{ since use min aggregator})$$



# Examples of linear constraints

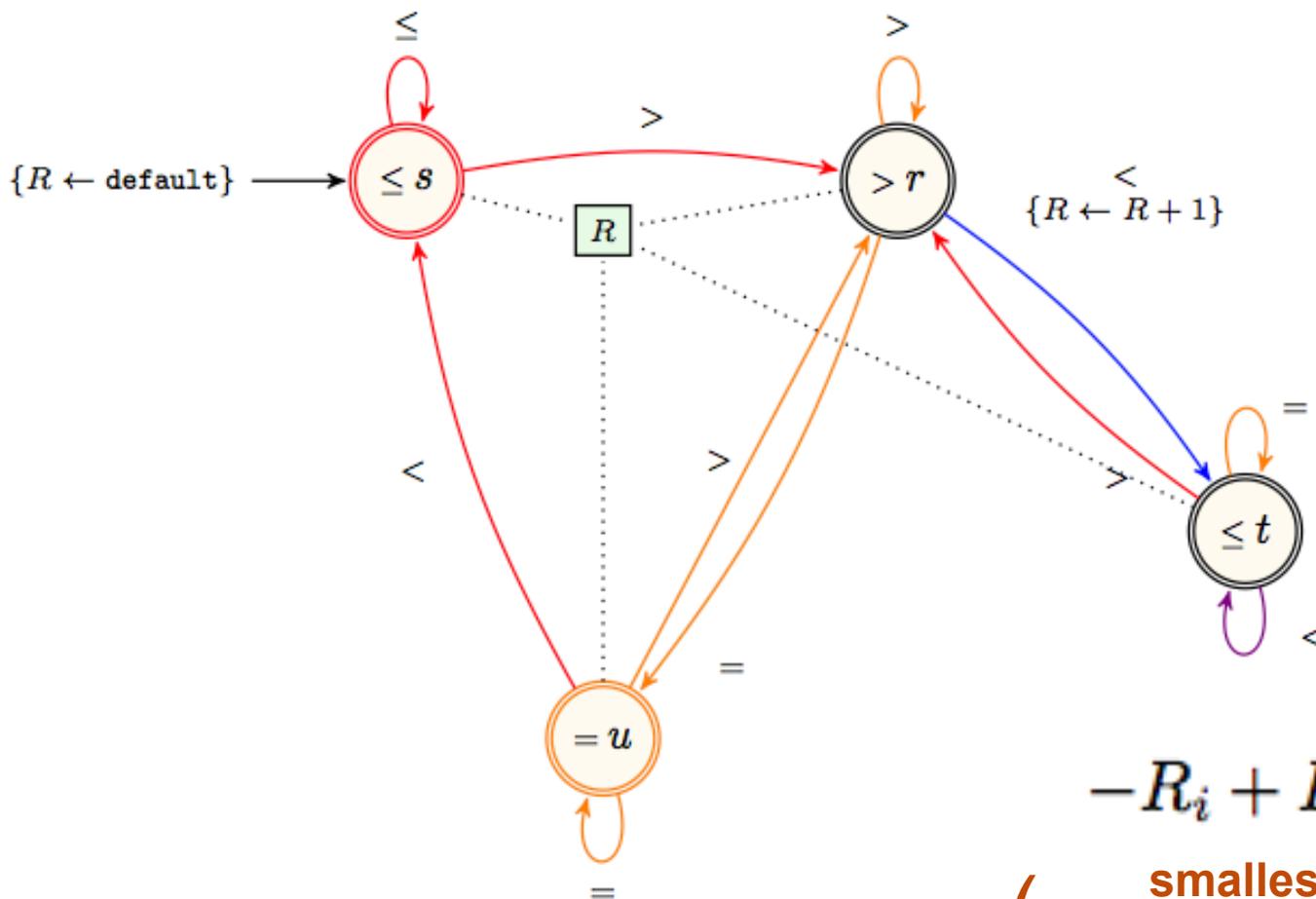
(nb\_dip\_on\_increasing\_sequence)



$$-R_i + R_{i-3} + 1 \geq 0$$

( smallest cycle between  
two consecutive incrementation )

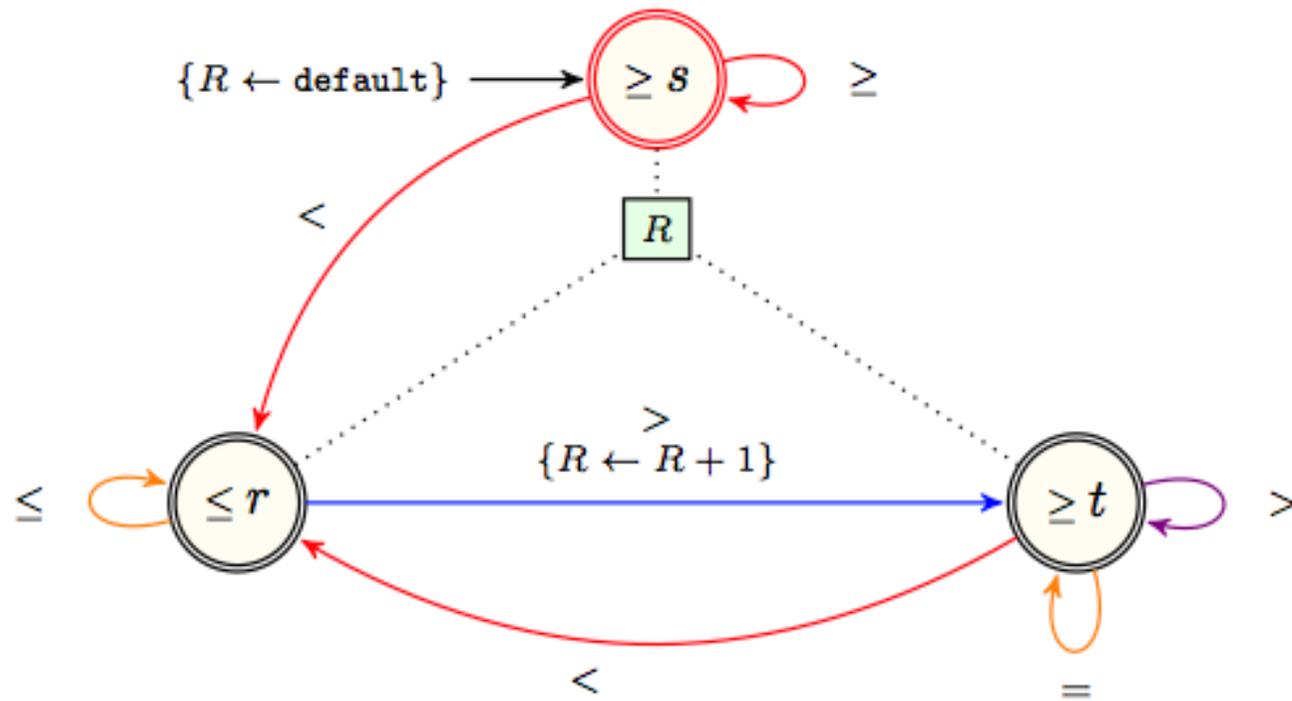
# Examples of linear constraints (nb\_gorge)



$$-R_i + R_{i-2} + 1 \geq 0$$

( smallest cycle between  
two consecutive incrementation )

# Examples of linear constraints (nb\_peak)

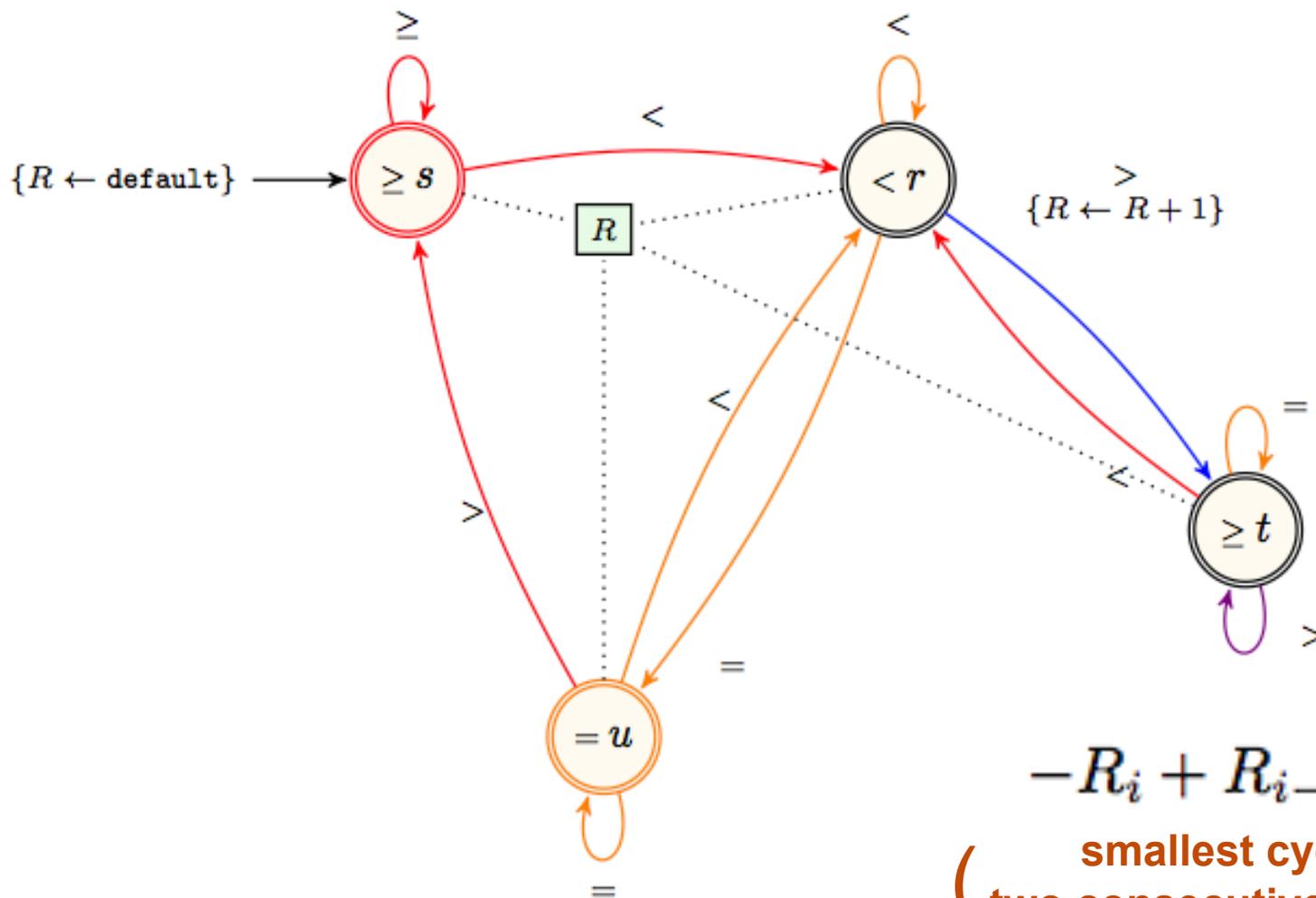


$$-R_i + R_{i-2} + 1 \geq 0$$

$$-2R_i + R_{i-1} + R_{i-2} + 2 \geq 0$$

( smallest cycle between  
two consecutive incrementation )

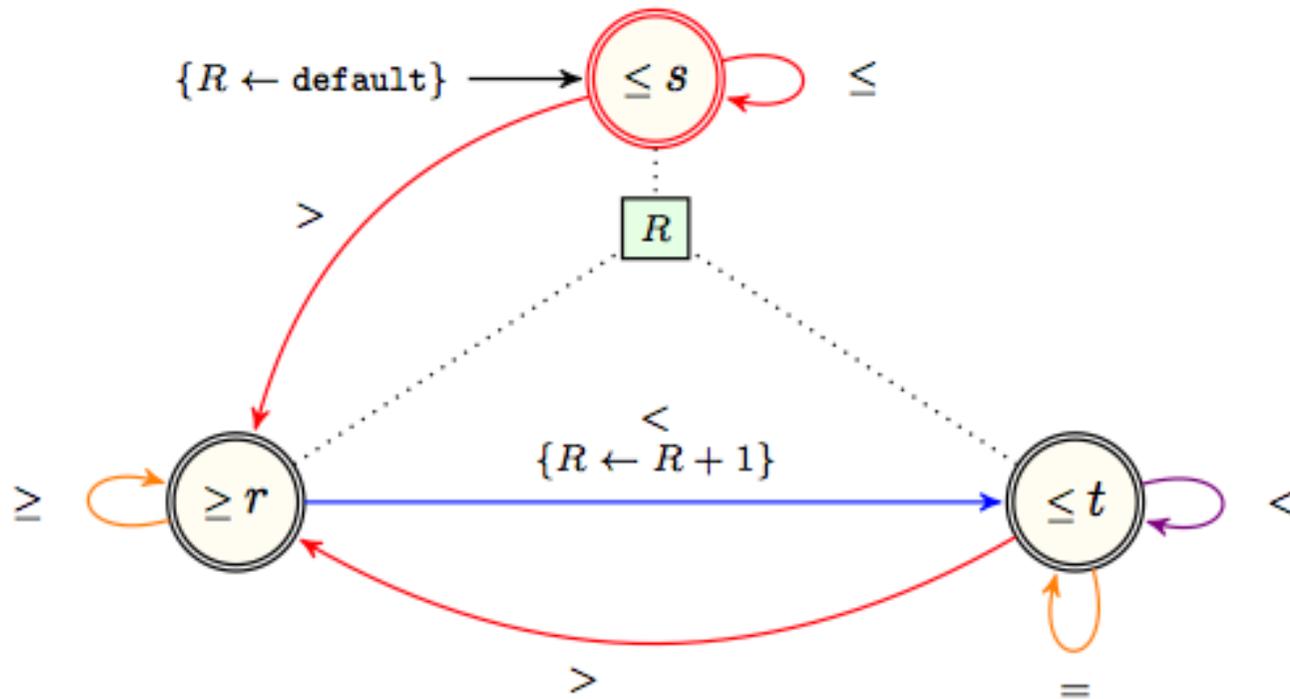
# Examples of linear constraints (nb\_summit)



$$-R_i + R_{i-2} + 1 \geq 0$$

( smallest cycle between  
two consecutive incrementation )

# Examples of linear constraints (nb\_valley)

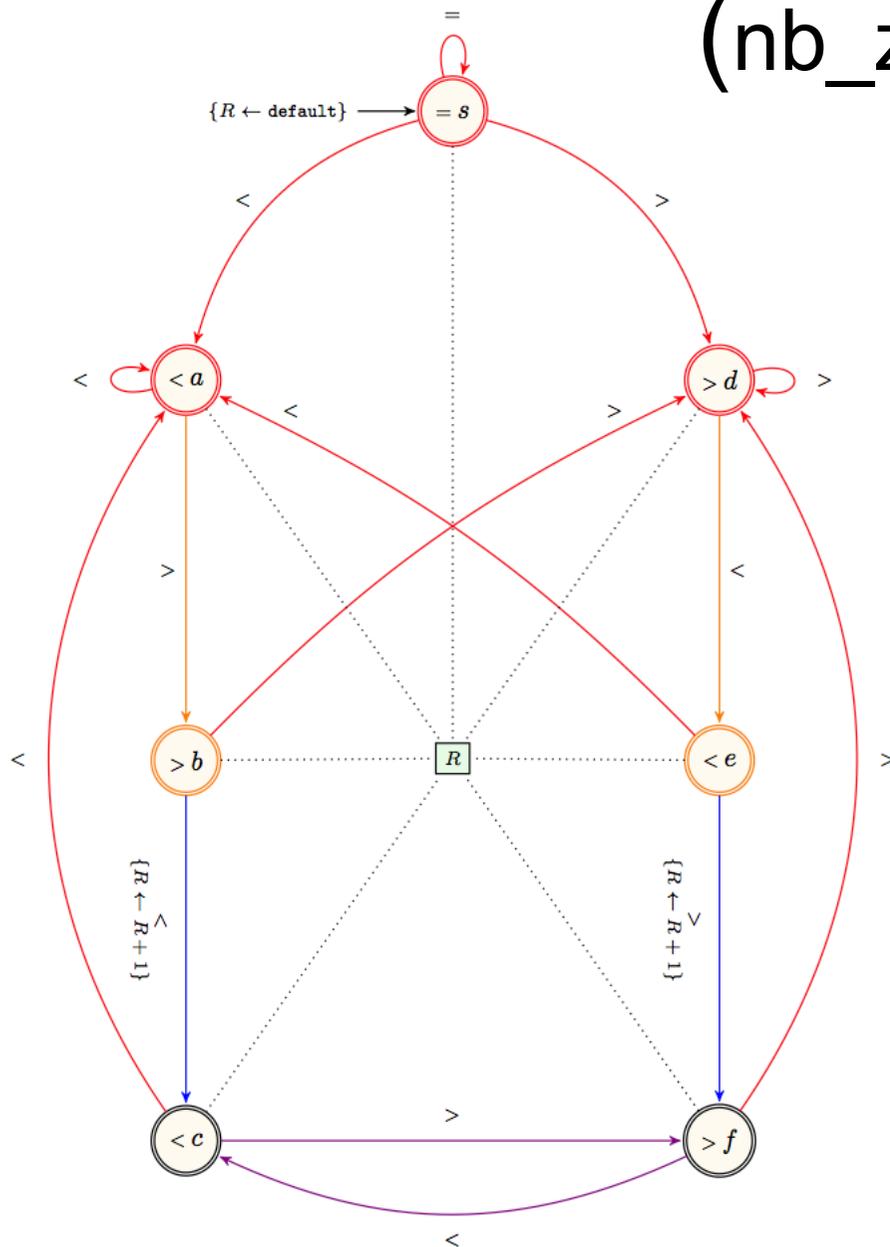


$$-R_i + R_{i-2} + 1 \geq 0$$

$$-2R_i + R_{i-1} + R_{i-2} + 2 \geq 0$$

( smallest cycle between  
two consecutive incrementation )

# Examples of linear constraints (nb\_zigzag)



$$R_i - R_{i-1} \geq 0$$

( increasing R since use  
sum aggregator and feature one )

$$-R_i + R_{i-3} + 1 \geq 0$$

( smallest cycle between  
two consecutive incrementation )

- Background
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- **Bounds**

# The need for bound

- Want to have lower/upper bound on the result returned by a time-series constraint parametrized by:
  - The sequence length
  - The smallest or largest values of the variables in the sequence

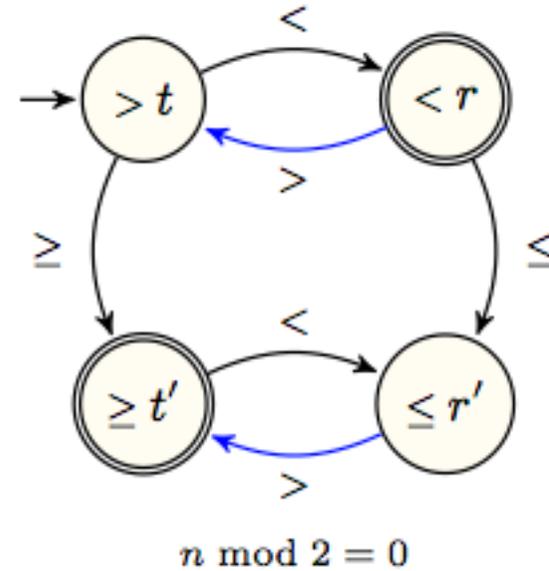
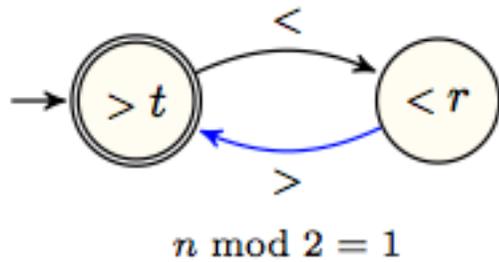
*Having a bound is good,  
but having a way to characterize  
all solutions reaching this bound is even better*

# Point

- In a significant number of cases we can **just use the standard regular constraint** (with an automaton having a **fixed number of states**) for characterizing all solutions reaching a given bound

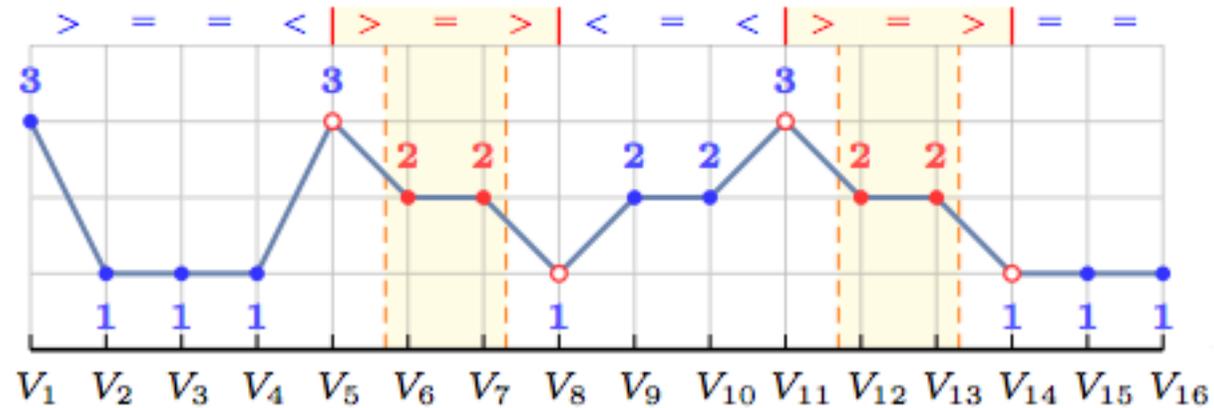
when the bound does not depend of the domain size

# Example 1: nb\_peak (upper bound $\lfloor \frac{n-1}{2} \rfloor$ )



0, 1, 0    0, 0, 1, 0    0, 1, 0, 0    0, 1, 0, 1    0, 1, 1, 0    1, 0, 1, 0  
 <>    =<>    <>=    <><    <=>    ><>

Exercise: get a transducer for the decreasing\_terrace pattern ( $\geq^+ \geq$ )



# Exercise: transducer for decreasing\_terrace

